

Duality and Geometry Straightness, Characterization and Envelope

Jean-Marc Chassery

LIS Grenoble

David Coeurjolly and Isabelle Sivignon

LIRIS Lyon

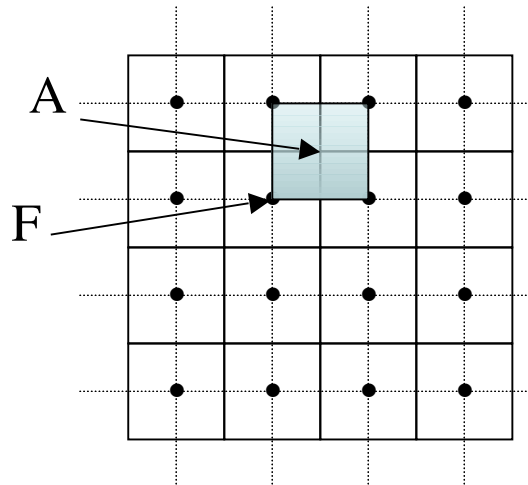
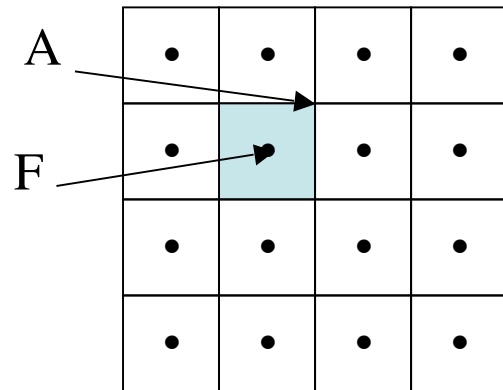


Outline

1. Introduction
2. Hough transform and discrete geometry
3. Preimage of a pixel, a set of pixels, a connected straight line
4. Preimage and Farey sequences
5. Extension in 3D
6. Duality versus polarity
7. Polygon enclosure and inclusion problems
8. Application to the convex approximation by interior.
9. Conclusion



1. Duality - Partitioning and meshes



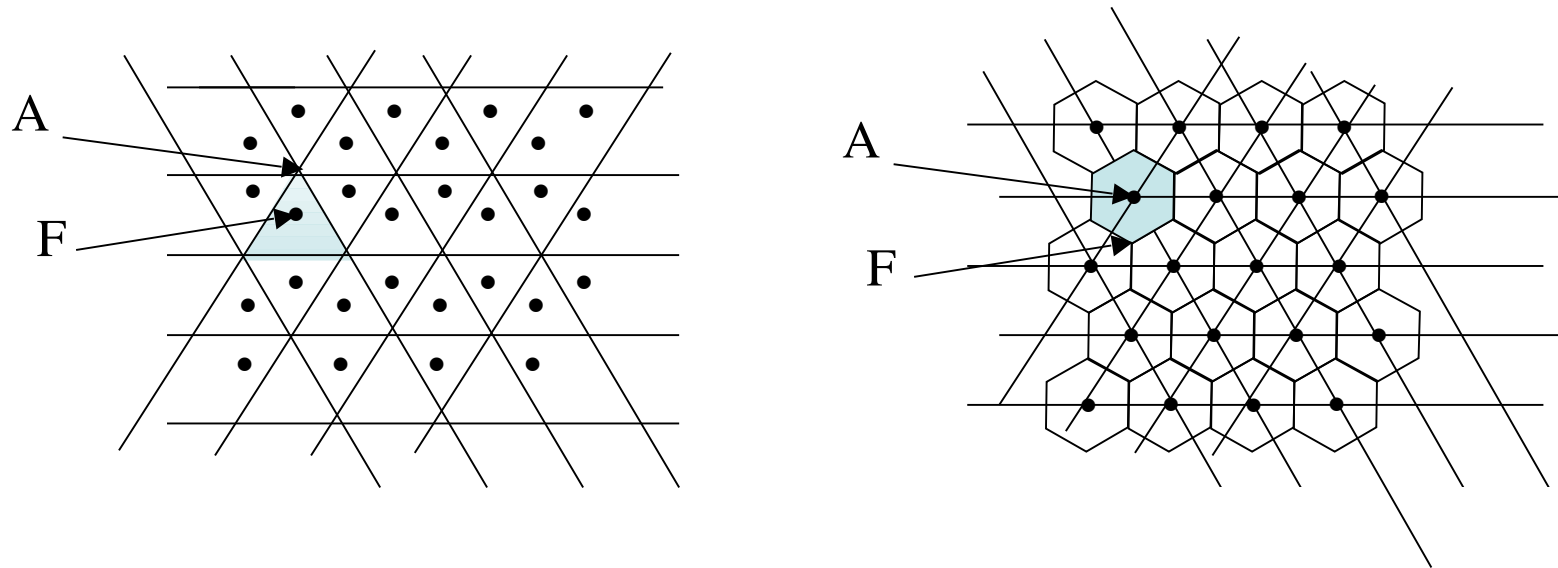
Let P be a partitioning.

If A is a top point of P and F a face of P with $A \subset F$ then
in the dual of P , A is a face and F a top point with $F \subset A$.

The dual of a square partitioning is a square partitioning.



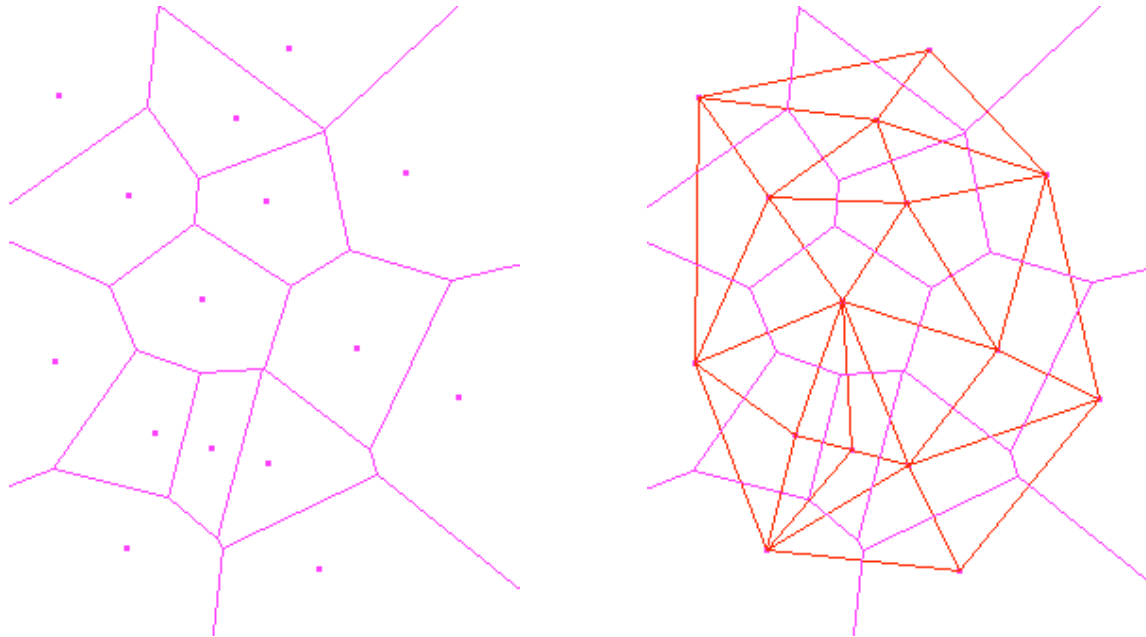
Duality - Partitioning and meshes



The dual of a triangular partitioning is an hexagonal partitioning.



Duality - Voronoi - Delaunay



The dual of a Voronoi partitioning is a Delaunay triangulation.

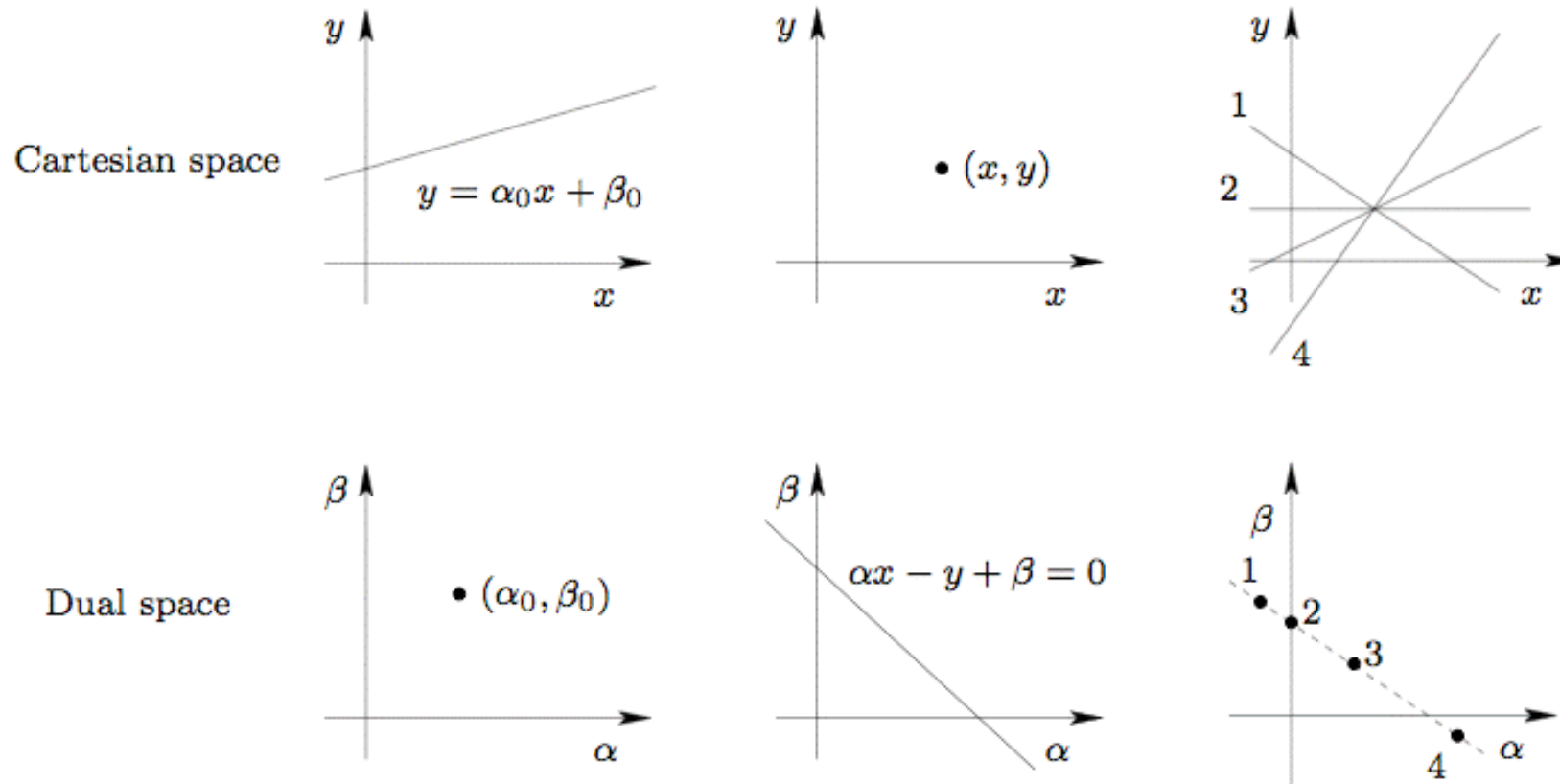


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2. Hough Transform



- A point in the (x,y) space matches up with one curve in the dual space
- A point in the dual space matches up with one curve in the (x,y) space.
- Points lying on a same line in (x,y) space match up with concurrent curves in the dual space.
- Points on a same curve in dual space match up with concurrent straight lines in (x,y) space.

Hough in discrete geometry

- Each point of parametric space can be seen either as a digital straight line or as an euclidean line, the discretization of which corresponding to the digital straight line.

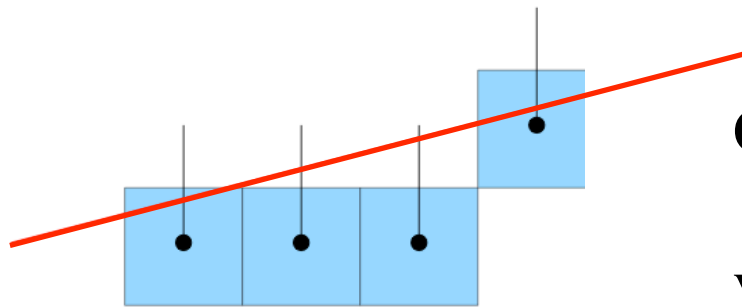
- In discrete geometry a naive straight line is not given by an equality but by 2 inequalities:

$$0 \leq \alpha x - y + \beta < 1$$

So a pixel corresponds to the intersection of 2 half spaces in (α, β) parameter space.



Digitization process for a 8-connected Digital straight line segment



Object Boundary Quantization (OBQ) :

$$y = \text{Floor} (\alpha x + \beta) \text{ with } x \in \mathbb{Z}$$



From Pixel to strip in the dual space

Let $P(x_0, y_0)$ be a pixel and $L : y = \alpha x + \beta$ a straight line

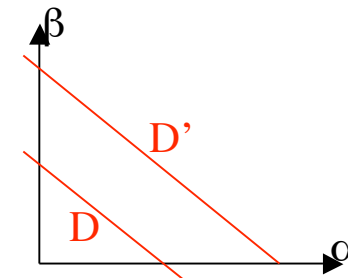
$$P \in \text{OBQ}(L) \Leftrightarrow y_0 = \text{Floor}(\alpha x_0 + \beta)$$

$$0 \leq \alpha x_0 + \beta - y_0 < 1$$

Strip noted by $B(x_0, y_0) = \{ \text{straight lines } L ; P \in \text{OBQ}(L) \}$

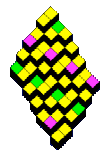
(α, β) such that:

$$B(x_0, y_0) = \left\{ \begin{array}{l} 0 \leq \alpha x_0 + \beta - y_0 \quad (\text{D}) \\ \text{and } \alpha x_0 + \beta - y_0 < 1 \quad (\text{D}') \end{array} \right\}$$



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Preimage of a set of pixels

Definition:

Let $S = \{(x_i, y_i)\}$ be a set of pixels, its preimage S_p is:

$$S_p = \{(\alpha, \beta) \mid 0 \leq \alpha x_i + \beta - y_i < 1 ; \forall i\}$$

$$S_p = \bigcap_i B(x_i, y_i)$$

$(\alpha, \beta) \in S_p \Leftrightarrow$ the straight line with parameters (α, β)
crosses the interval $[y_i, y_i+1[$ of each pixel.

Dorst & Smeulders, PAMI, 1984

McIlroy AT&T Tech Journal, 1985

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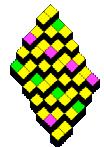
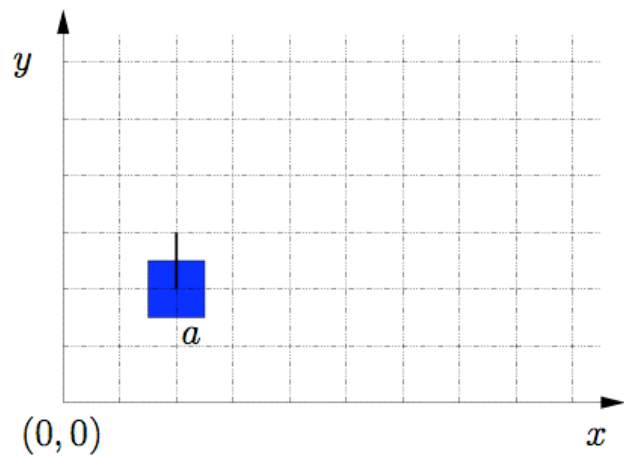
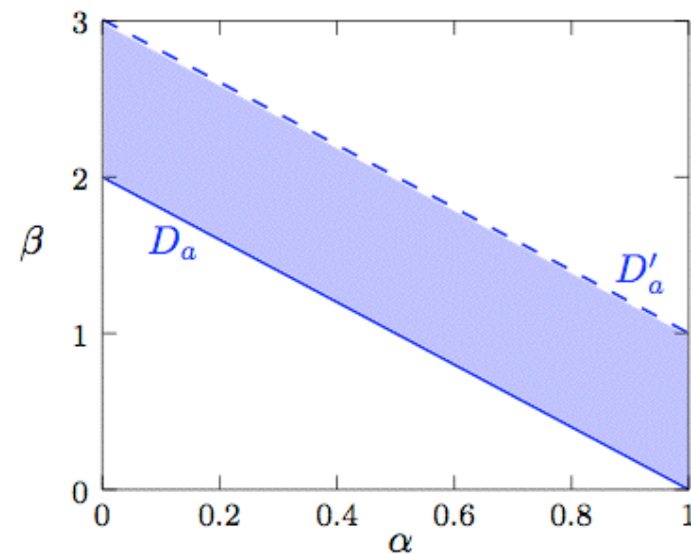


Illustration of preimage



$$a=(2,2)$$

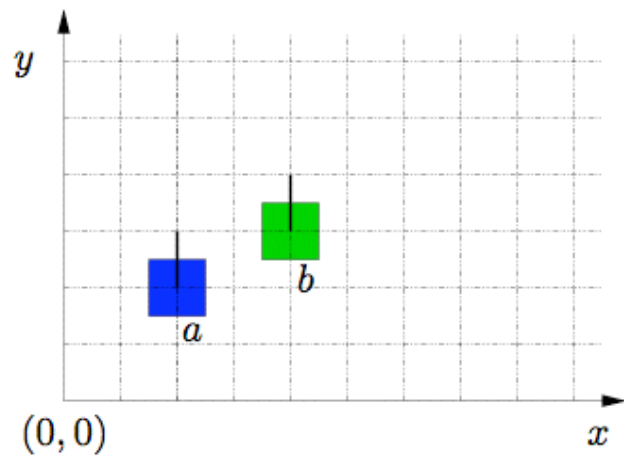


$$D_a : 2\alpha + \beta - 2 = 0$$

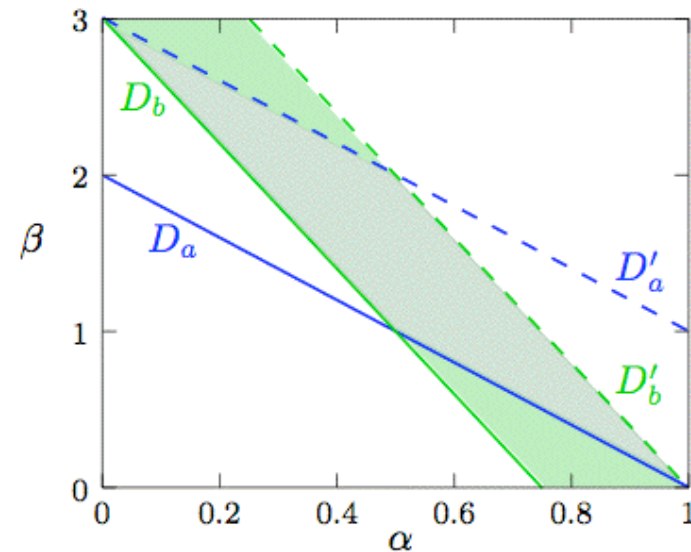
$$D'_a : 2\alpha + \beta - 2 = 1$$



Illustration of preimage



$b=(4,3)$

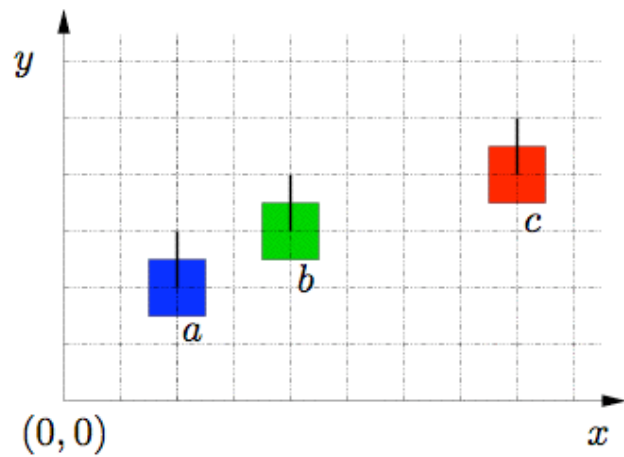


$$D_b : 4\alpha + \beta - 3 = 0$$

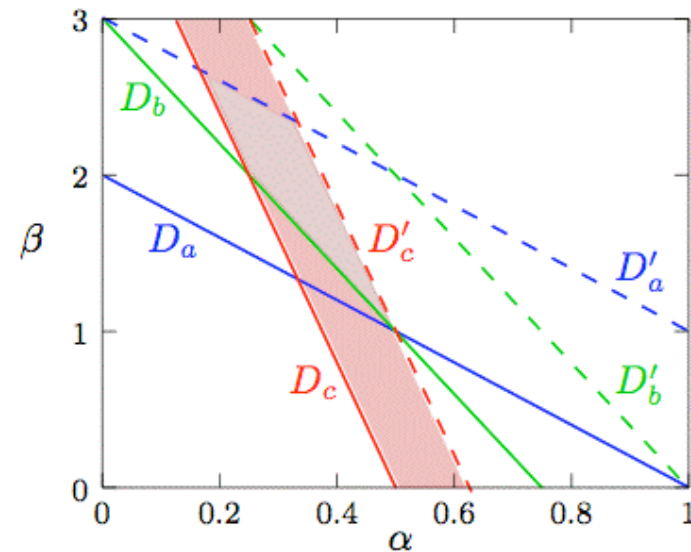
$$D'_b : 4\alpha + \beta - 3 = 1$$



Illustration of preimage



$$c=(8,4)$$



$$D_c : 8 \alpha + \beta - 4 = 0$$

$$D'_c : 8 \alpha + \beta - 4 = 1$$

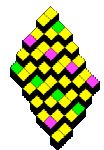
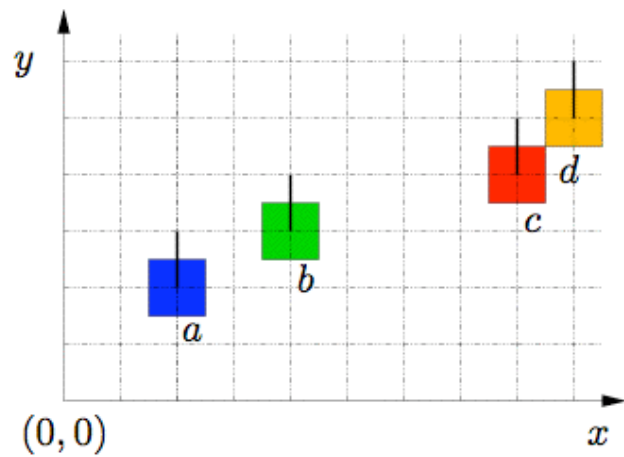
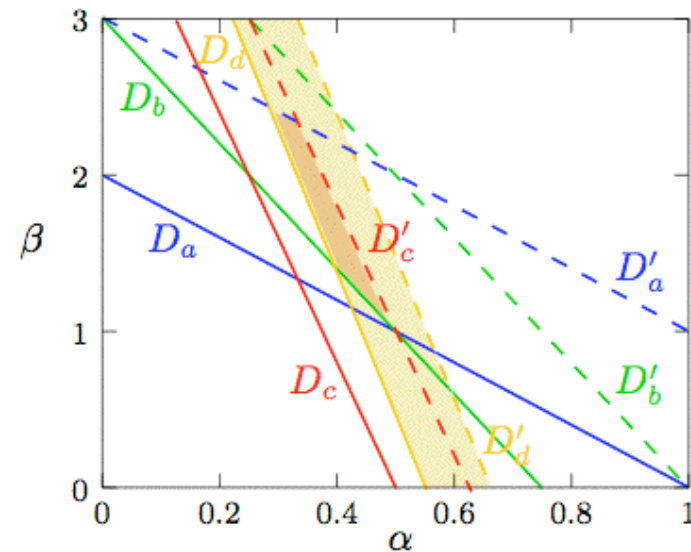


Illustration of preimage



$d=(9,5)$



$$D_d : 9 \alpha + \beta - 5 = 0$$

$$D'_d : 9 \alpha + \beta - 5 = 1$$

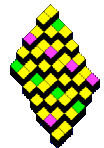
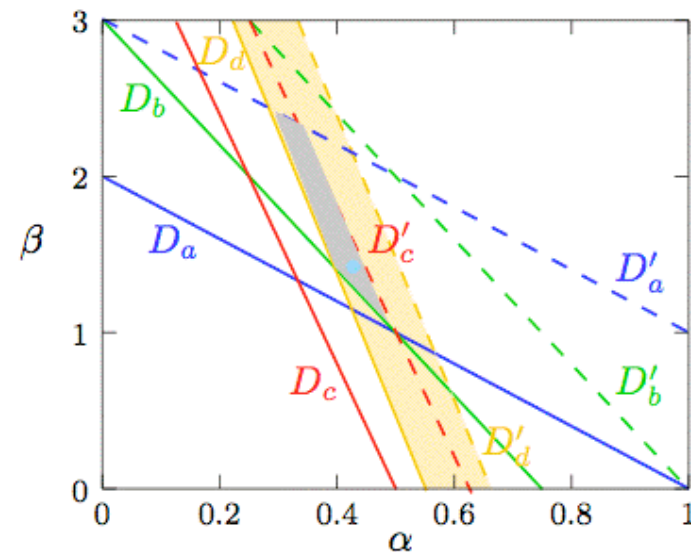
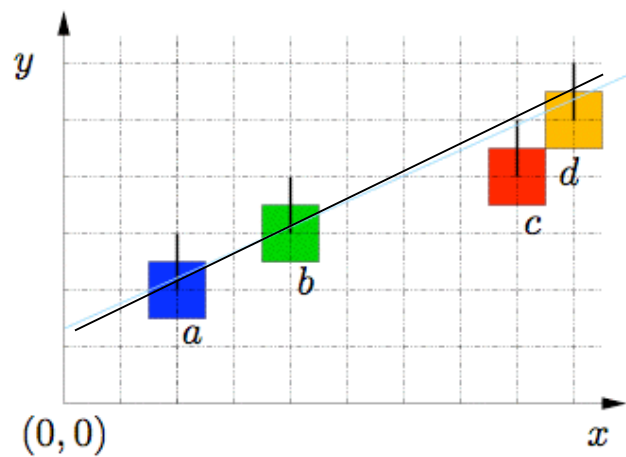


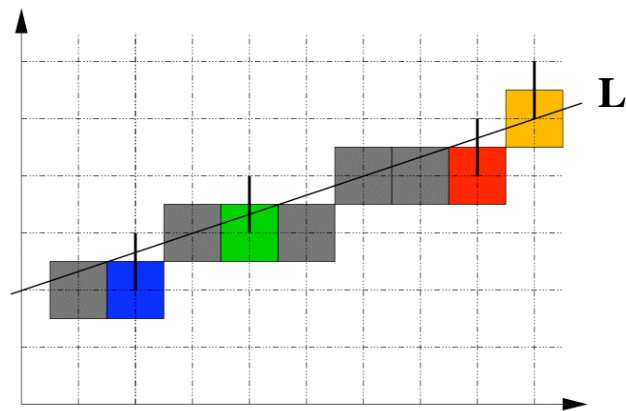
Illustration of preimage



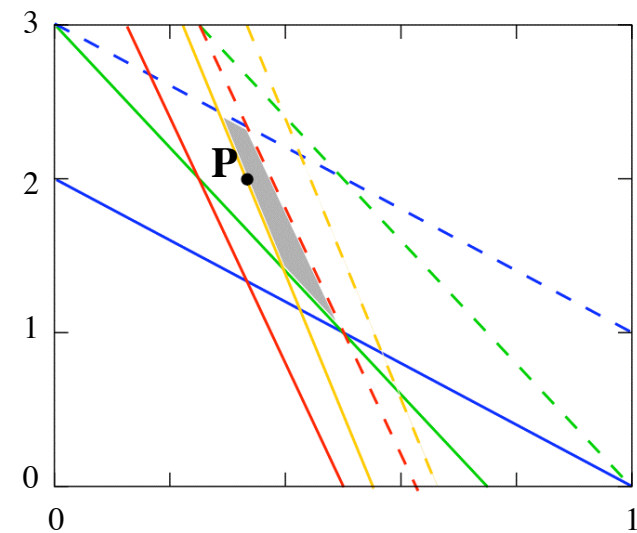
Set of solutions



Illustration of preimage



$$L : \frac{1}{3}x - y + 2 = 0$$

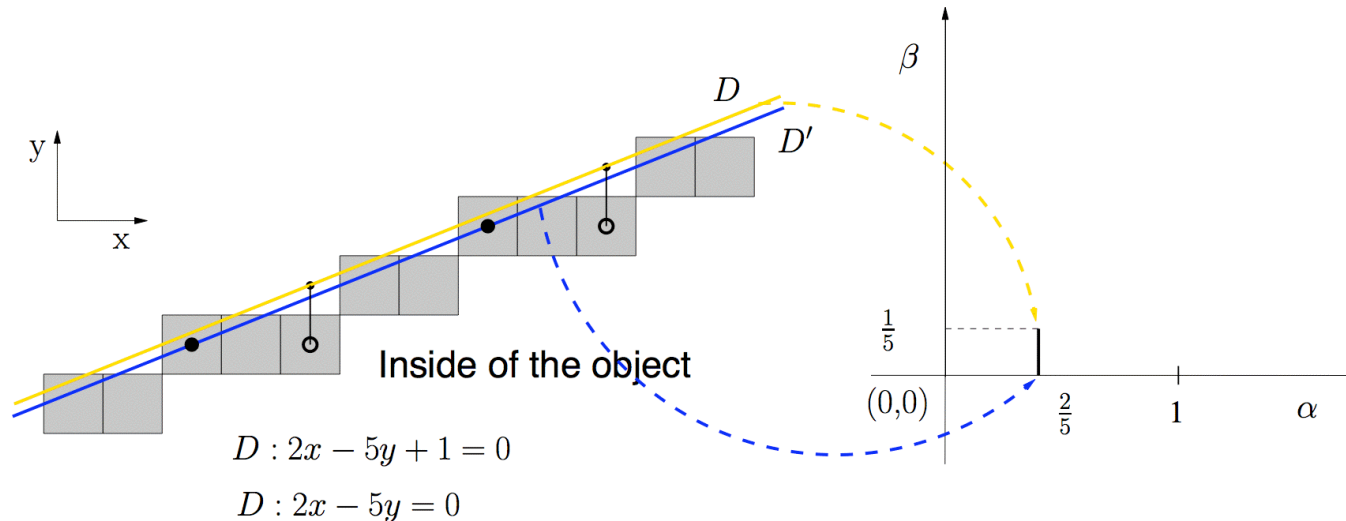


P (1/3, 2) involves the straight line L



Preimage of a straight line

In case of a straight line with parameters (α_0, β_0) , the Preimage of the discretized infinite set (the naive line) is a vertical segment with extremities $(\alpha_0, 0)$ and (α_0, β_0) .



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Farey series

Definition: The Farey series F_m of order m is the ascending series of irreducible fractions h/k between 0 and 1 whose denominators do not exceed m .

Property

- if h/k and h'/k' are two successive terms in F_m (with $h/k < h'/k'$), then $kh' - hk' = 1$.
- if h/k , h''/k'' and h'/k' are 3 successive terms in F_m (with $h/k < h''/k'' < h'/k'$), then $h''/k'' = (h+h')/(k+k')$. The fraction h''/k'' is called the **mediant** of h/k and h'/k' .
- F_{m+1} can be computed from F_m adding the mediant with denominator less or equal than $m + 1$ of each two successive fractions in F_m

F_1	{	0/1,				1/1 }
F_2	{	0/1,		1/2,		1/1 }
F_3	{	0/1,		1/3,	1/2,	2/3, 1/1 }
F_4	{	0/1,		1/4, 1/3, 2/5, 1/2, 3/5, 2/3, 3/4,		1/1 }

Farey Fan

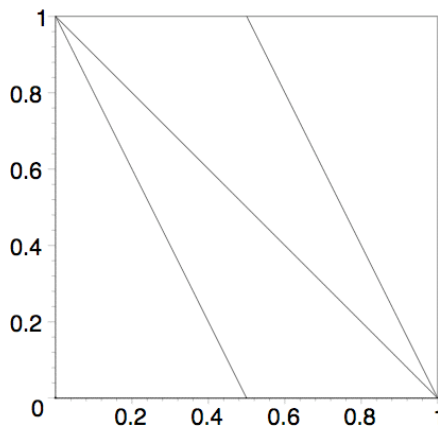
Definition (Rays)

$$R(x, y) = \{(\alpha, \beta) \mid \beta = -\alpha x + y\}$$

According to the strip decomposition, for each point (x, y) , we have the pair of parallel rays $R(x, y)$ and $R(x, y + 1)$ in the (α, β) dual space.

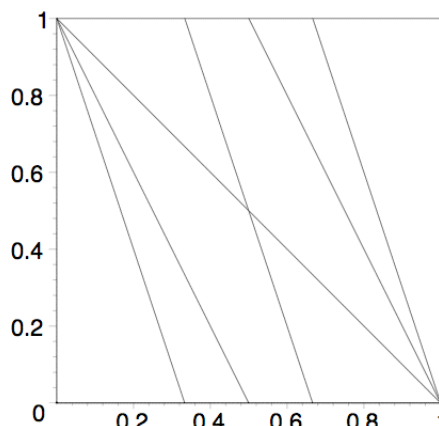
Farey Fans

If we focus on the domain $[0, 1] \times [0, 1[$ and segments in the first octant, rays $R(x, y)$ with $0 \leq y \leq x \leq n$ define a diagram called the Farey fan of order n .



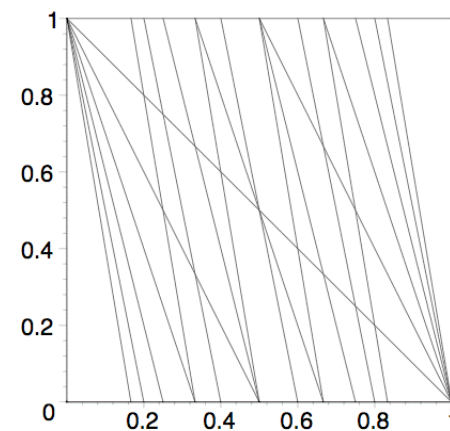
Farey series of order 2

$\{ 0/1, 1/2, 1/1 \}$



Farey series of order 3

$\{ 0/1, 1/3, 1/2, 2/3, 1/1 \}$



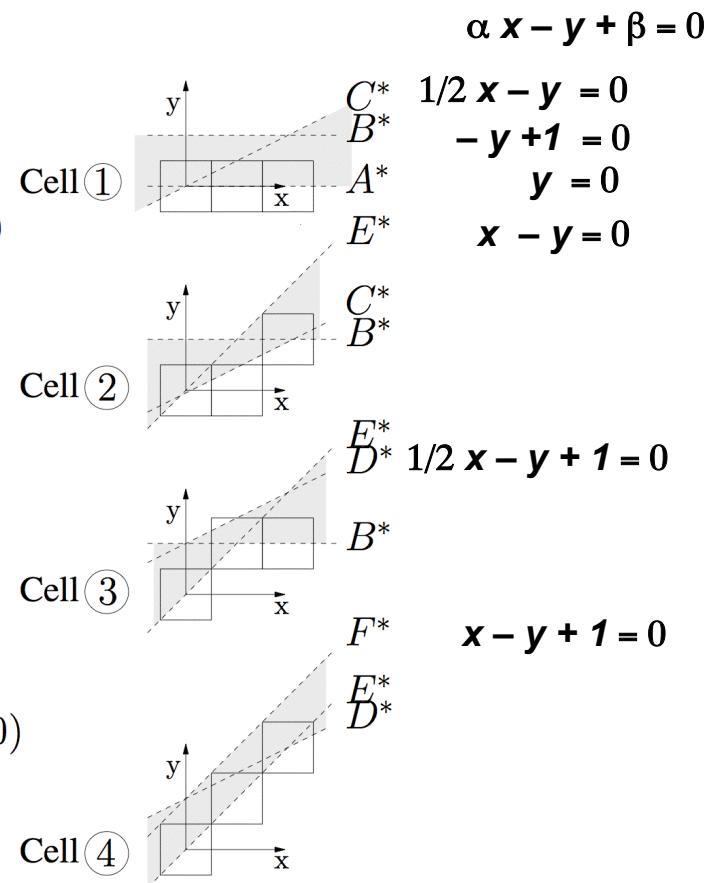
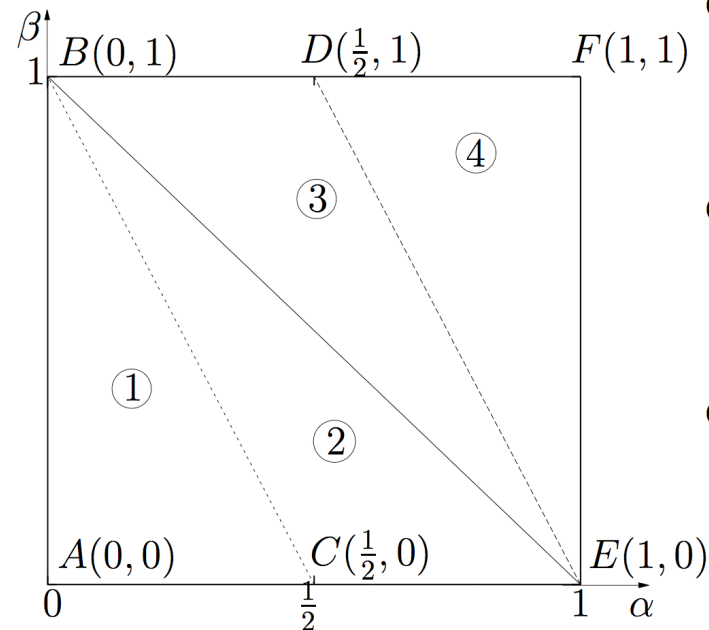
Farey series of order 6

$\{ 0/1, 1/6, 1/5, 1/4, 1/3, 2/5, 1/2, 3/5, 2/3, 3/4, 4/5, 5/6, 1/1 \}$

Farey fan and Digital Straight line

Farey fan of order n represents all the Digital Straight line of length $n+1$

Example: Farey fan of order 2

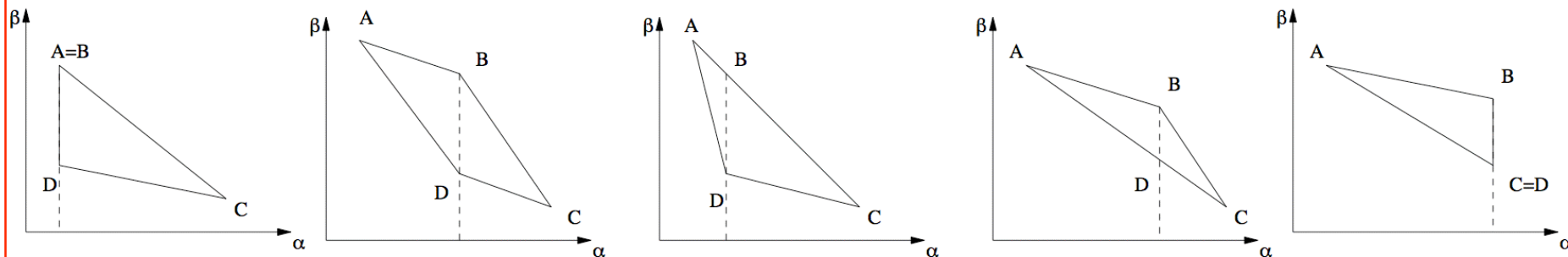


Preimage of a naive segment

Property:

Let S be a set of $(N+1)$ 8-connected pixels and x_0 be the minimum abscissa of this set of pixels. Then the preimage S_p of S has the following properties:

1. S_p is a convex polygon with at most four vertices;
2. Two consecutive abscissa of the vertices are consecutive terms in the Farey series of order $\max(x_0, N-x_0)$. Moreover, for a given abscissa equal to p/q , the corresponding ordinate is a multiple of $1/q$;
3. If this polygon has four vertices, then two out of the four vertices have the same abscissa.



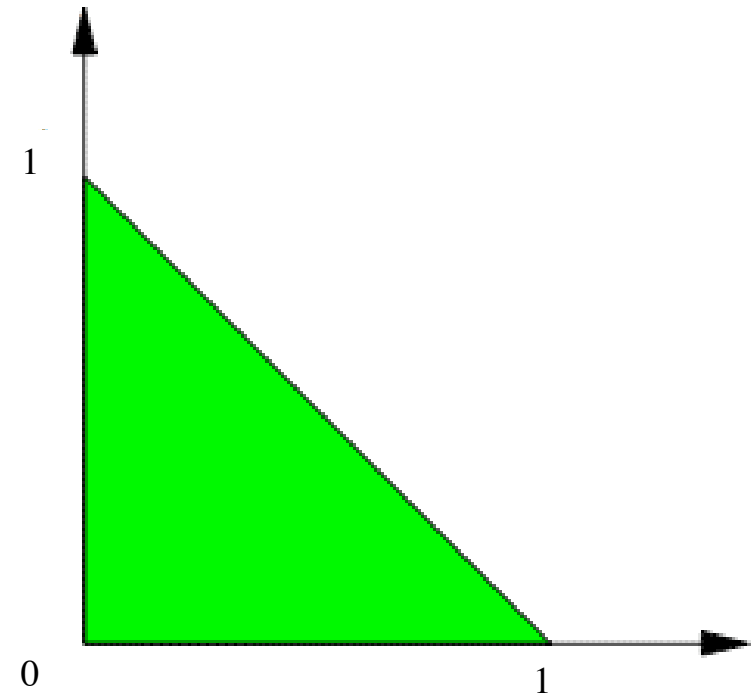
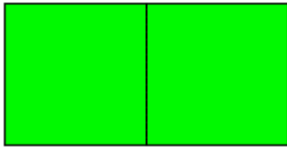
Recognition of a digital straight line based on preimage

Let S be a set of pixels.

Property: S is a piece of digital straight line iff
the preimage S_p of S is not empty.



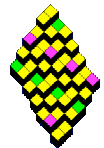
Illustration

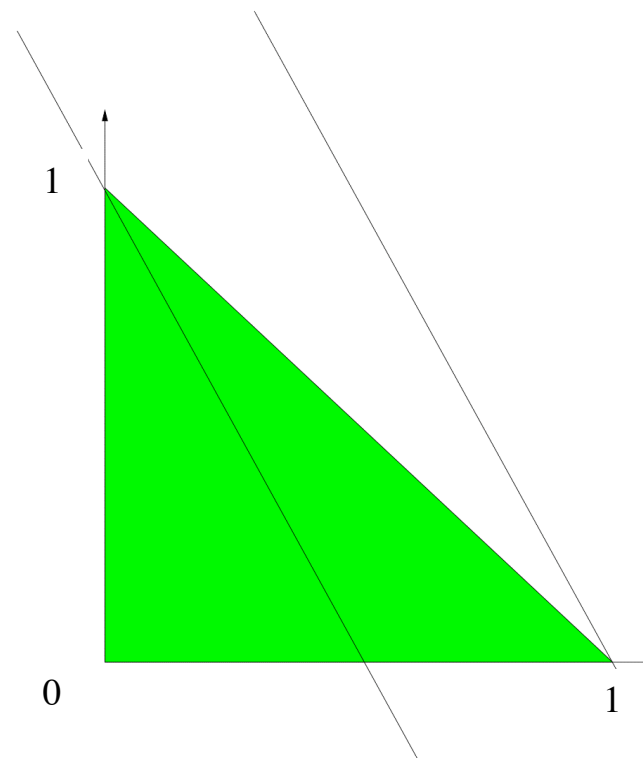
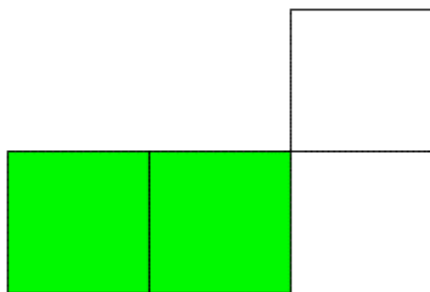


Step: Initialization

S_p coordinates: $\{(0, 1), (0, 0), (1, 0), (0, 1)\}$

slopes: $\{(_ _, (0, 0), (1, 1), (_ _)\}$



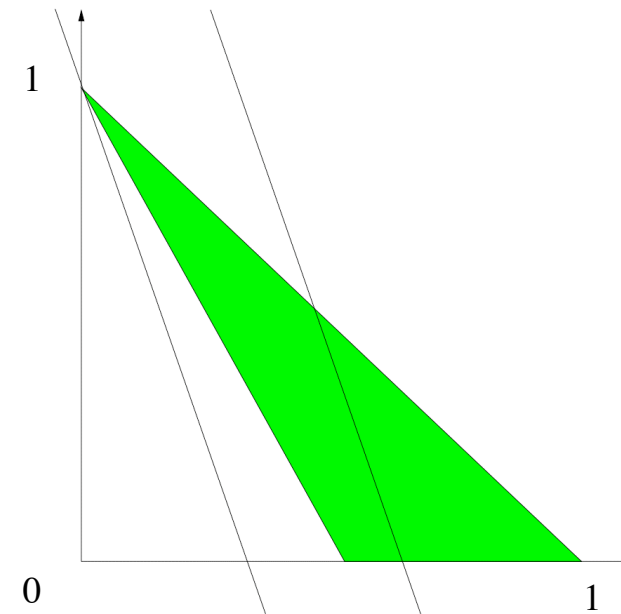
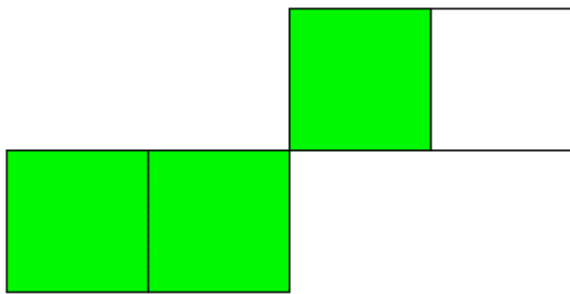


Step: $D \cap Llow \neq \emptyset$

S_p coordinates: $\{(0, 1), (1/2, 0), (1, 0), (1/2, 1/2)\}$

slopes: $\{(2, 1), (0, 0), (1, 1), (1, 1)\}$



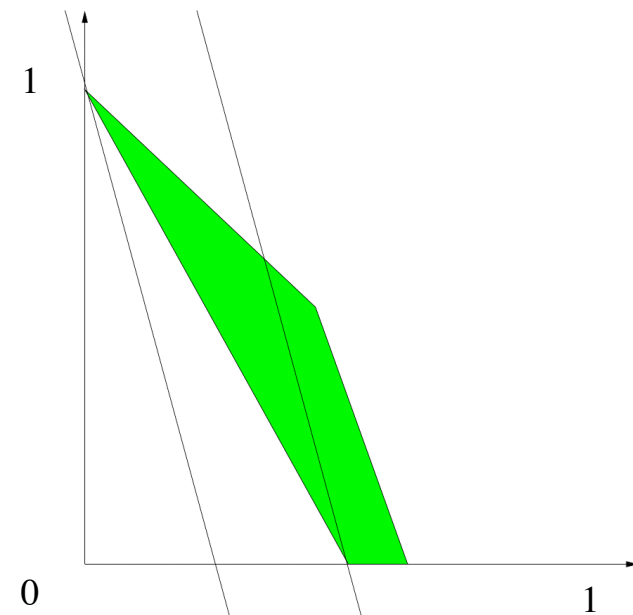
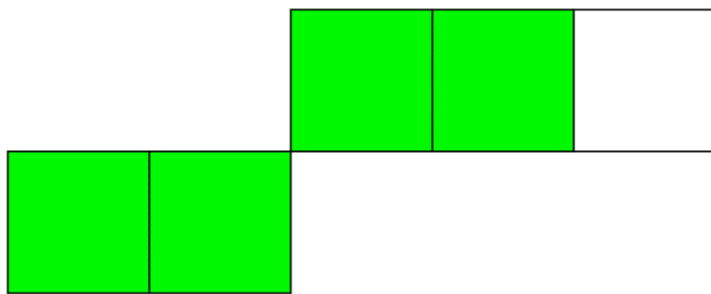


Step: $D \cap Lhigh \neq \emptyset$

S_p coordinates: $\{(0, 1), (1/2, 0), (2/3, 0), (1/2, 1/2)\}$

slopes: $\{(2, 1), (0, 0), (3, 1), (1, 1)\}$

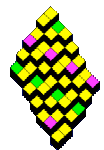


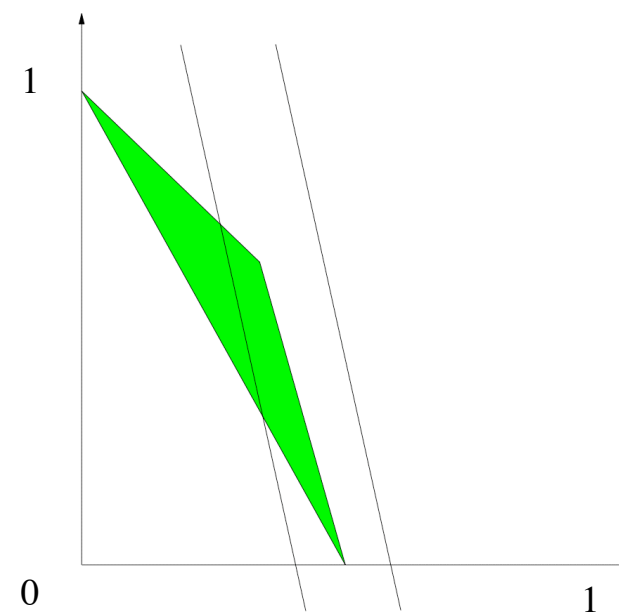
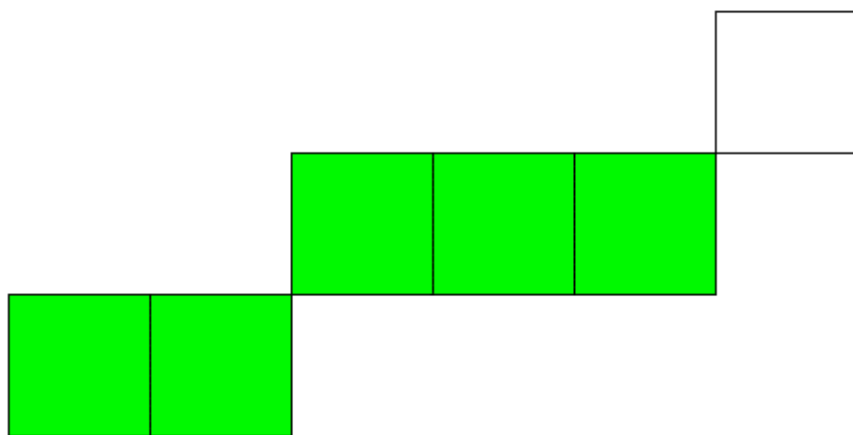


Step: $B \cap Lhigh \# \emptyset$

S_p coordinates: $\{(0, 1), (1/3, 1/3), (1/2, 0), (1/3, 2/3)\}$

slopes: $\{(2, 1), (0, 0), (4, 1), (1, 1)\}$



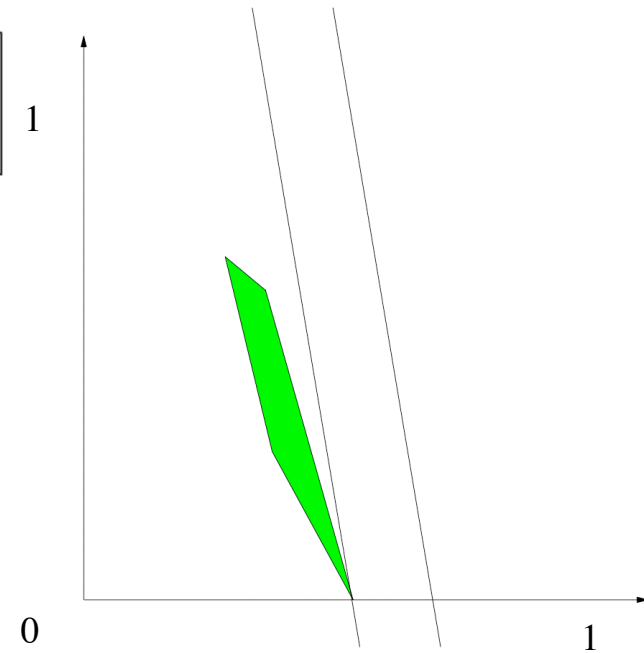
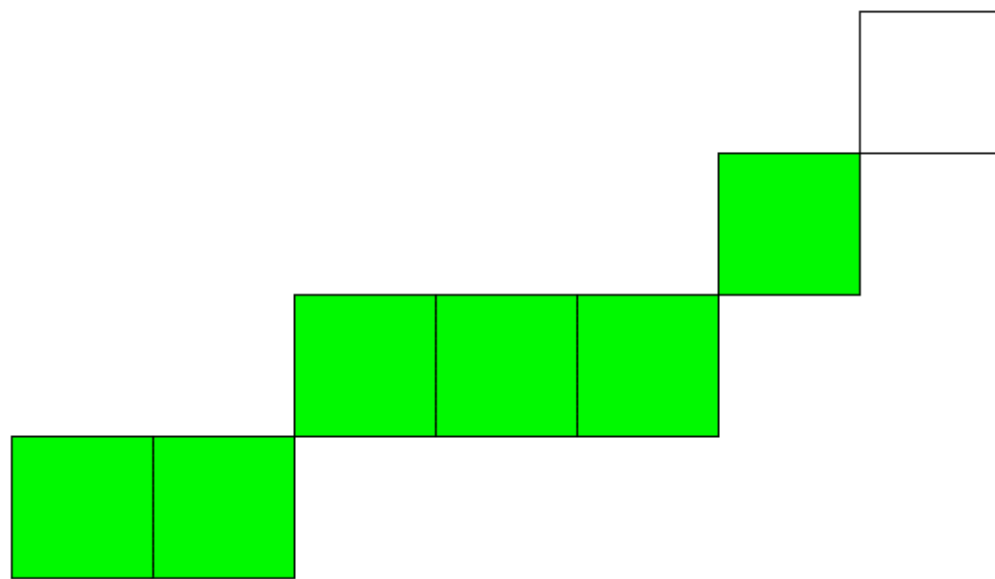


Step: $B \cap Llow \# \emptyset$

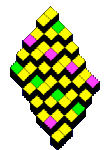
S_p coordinates: $\{(1/4, 3/4), (1/3, 1/3), (1/2, 0), (1/3, 2/3)\}$

slopes: $\{(5, 2), (0, 0), (4, 1), (1, 1)\}$



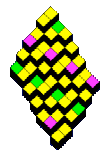


Step: $B \notin B(p)$ or $D \notin B(p)$
 S_p coordinates: $\emptyset$
slopes: $\emptyset$



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Digital Plane: Parametric Equation

$$0 \leq ax + by + cz + r < \omega$$

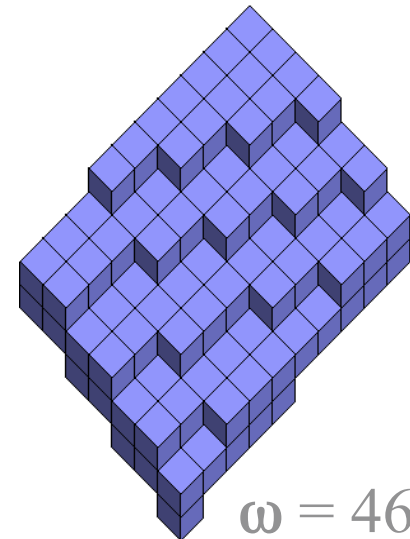
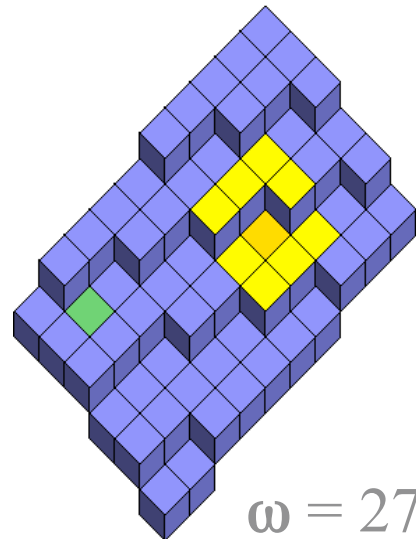
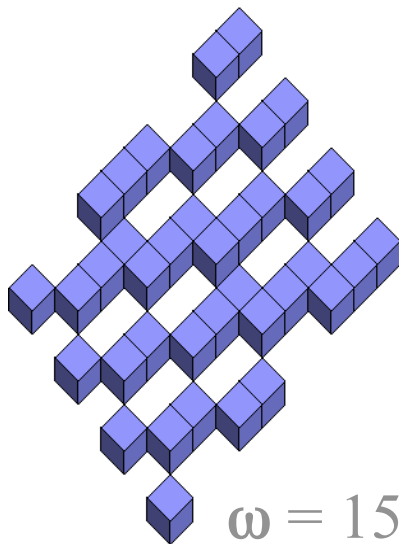
ω : *thickness*

(a,b,c) : *normal vector*

r : *translation parameter*

Example (a,b,c,r) :

$$P = (6, 13, 27, 0)$$



Digital Naive Plane

$$0 \leq ax + by + cz + r < \omega$$

ω : *thickness*

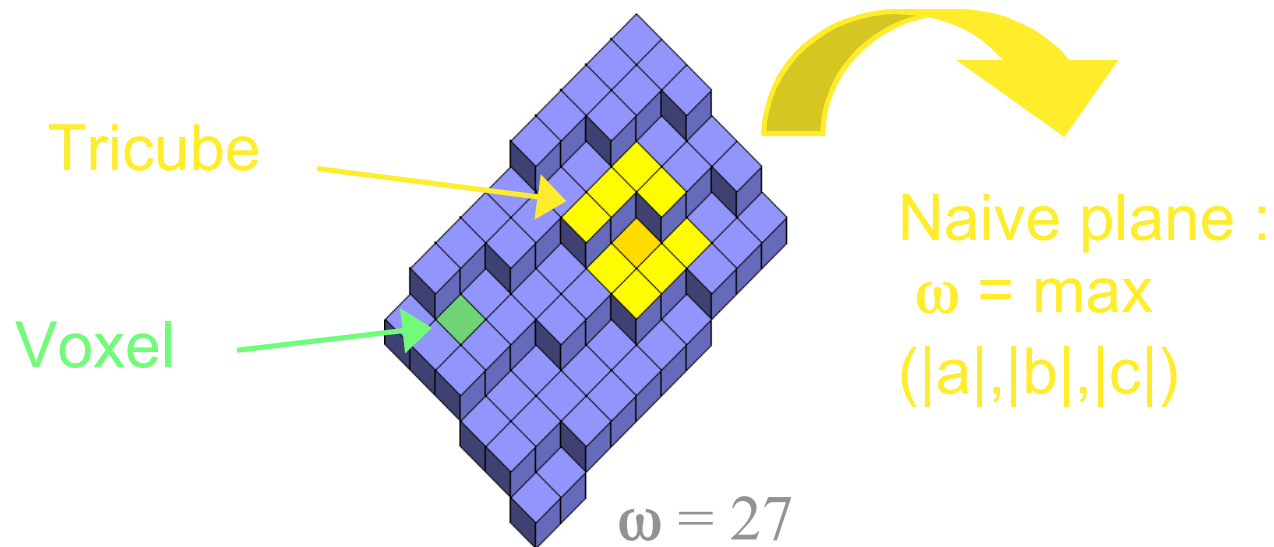
(a,b,c) : *normal vector*

r : *translation parameter*

Example

(a,b,c,r) :

$P = (6,13,27,0)$



Extension in 3D in the dual space

Let $P(x_0, y_0, z_0)$ be a voxel and $L : z = \alpha x + \beta y + \gamma$ a straight line

$$P \in \text{OBQ}(L) \Leftrightarrow z_0 = \text{Floor}(\alpha x_0 + \beta y_0 + \gamma)$$

$$0 \leq \alpha x_0 + \beta y_0 + \gamma - z_0 < 1$$

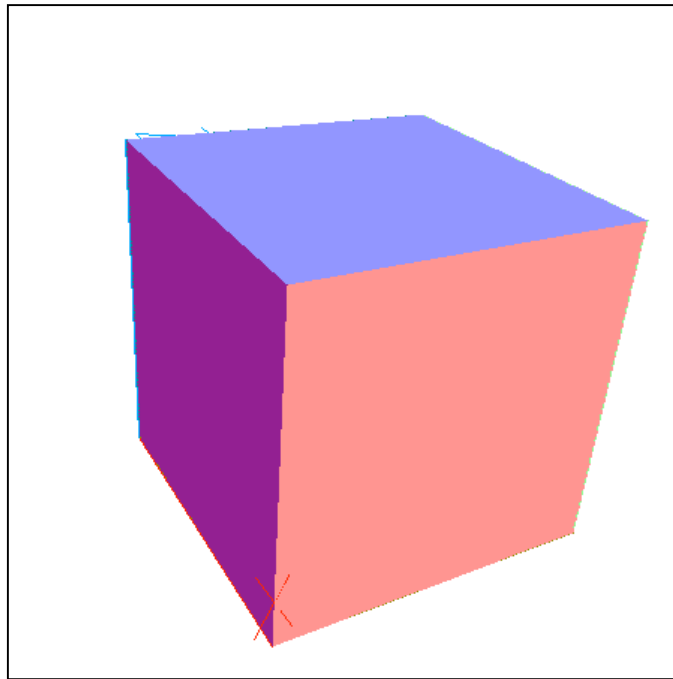
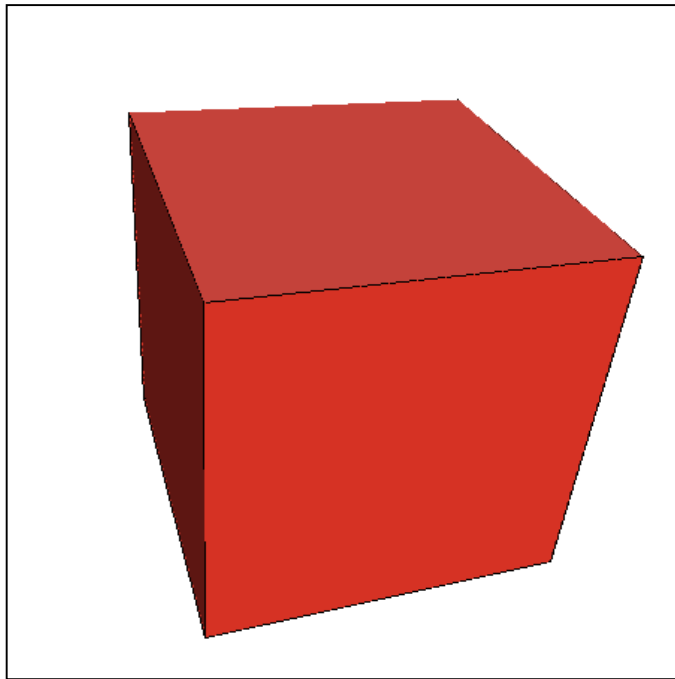
Strip associated to a voxel

(α, β, γ) such that:

$$B(x_0, y_0, z_0) = \left\{ \begin{array}{l} 0 \leq \alpha x_0 + \beta y_0 + \gamma - z_0 \text{ (D)} \\ \text{and } \alpha x_0 + \beta y_0 + \gamma - z_0 < 1 \text{ (D')} \end{array} \right\}$$



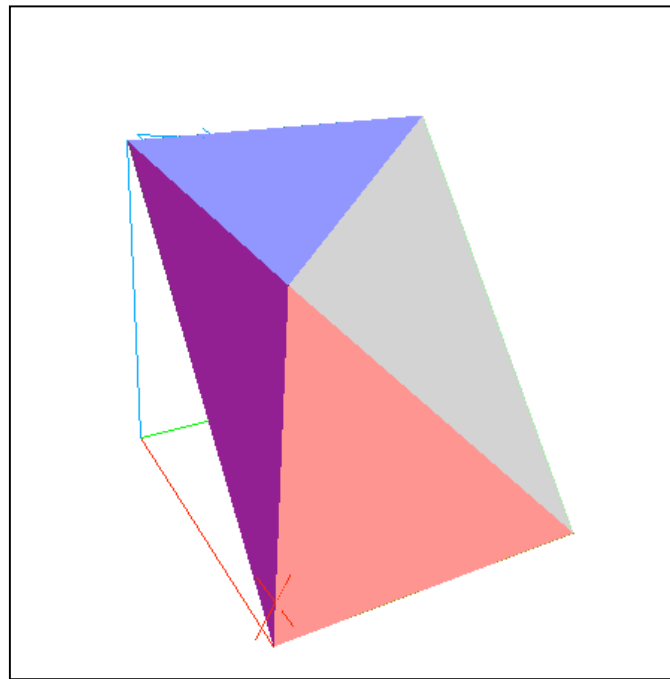
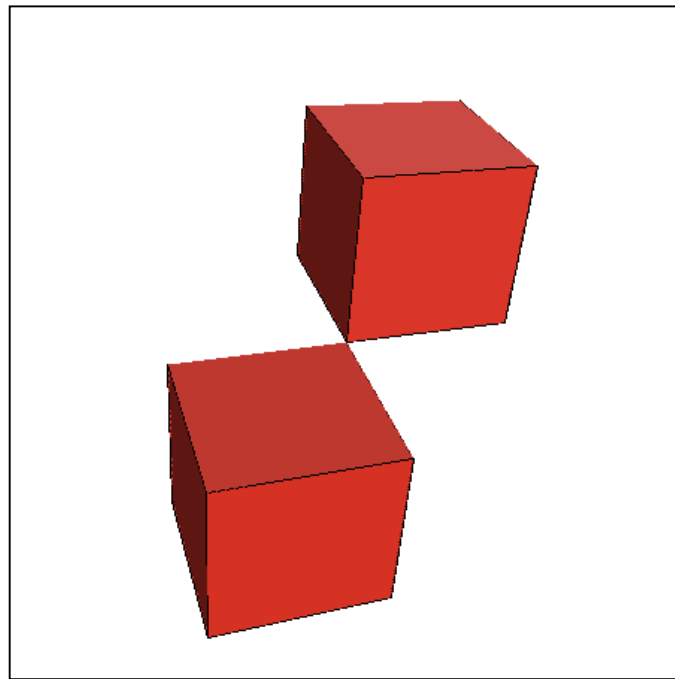
Recognition example



- $(0, 0, 0)$
- $(1, 1, 1)$
- $(2, 1, 1)$
- $(2, 2, 1)$
- $(3, 2, 1)$



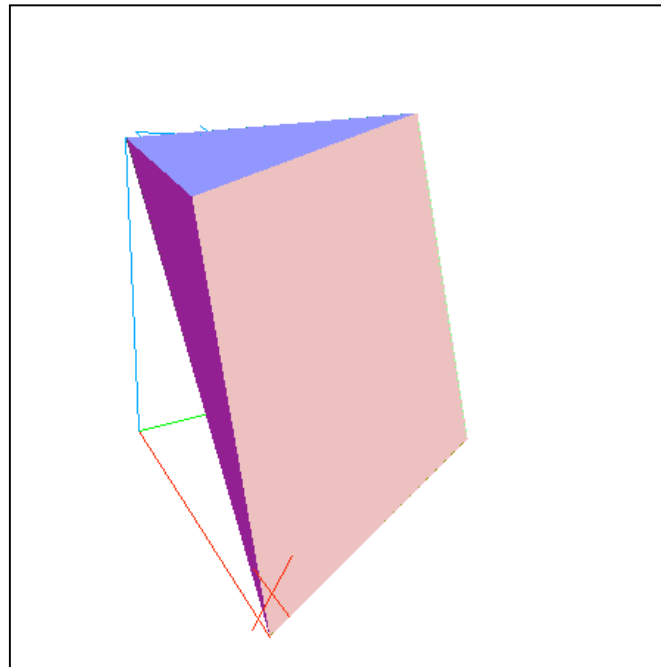
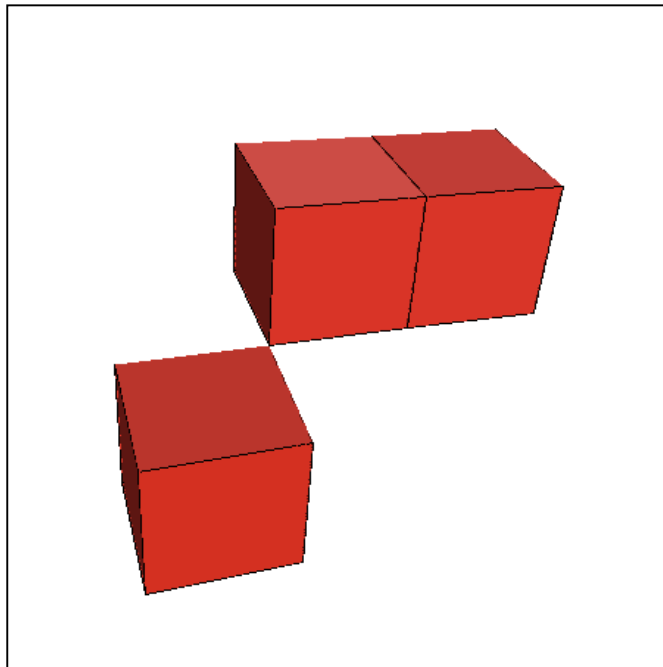
Recognition example



- $(0, 0, 0)$
- $(1, 1, 1)$
- $(2, 1, 1)$
- $(2, 2, 1)$
- $(3, 2, 1)$



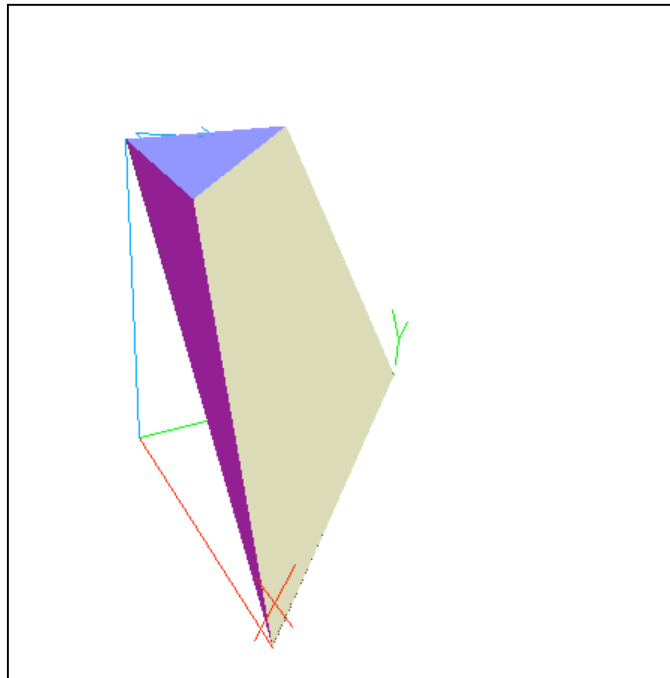
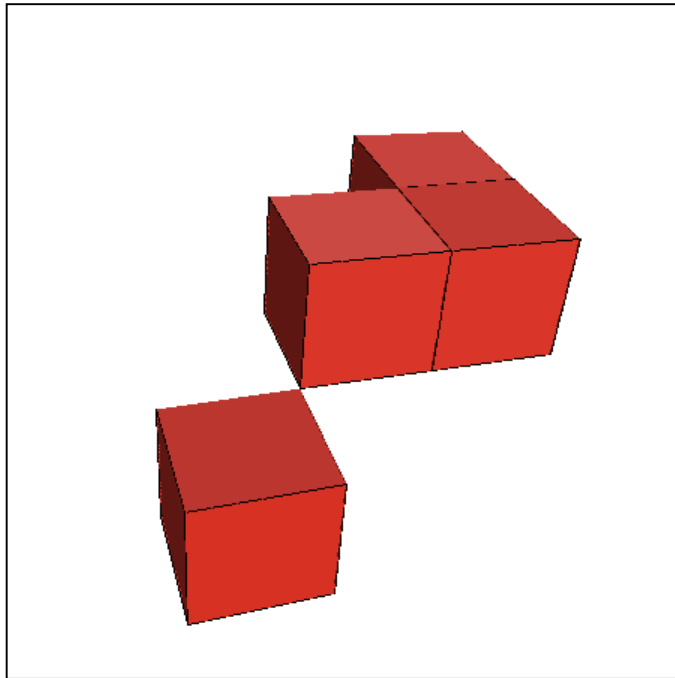
Recognition example



- (0, 0, 0)
- (1, 1, 1)
- (2, 1, 1)
- (2, 2, 1)
- (3, 2, 1)



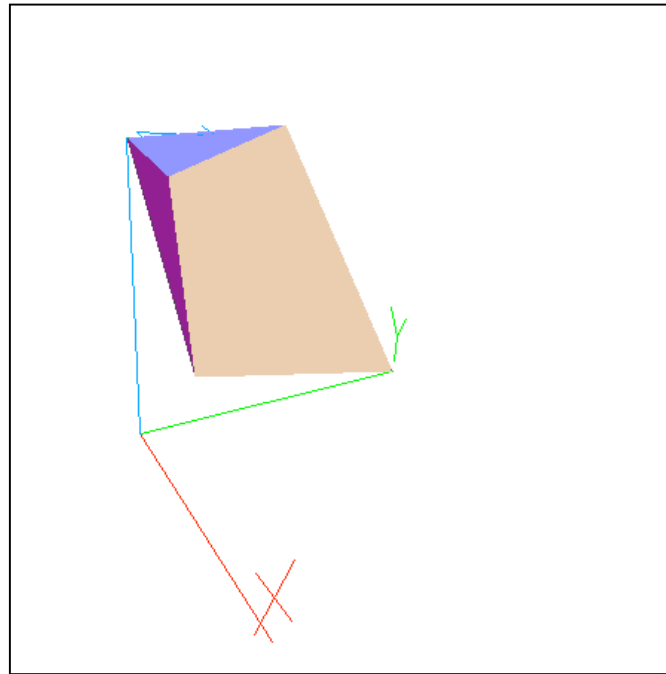
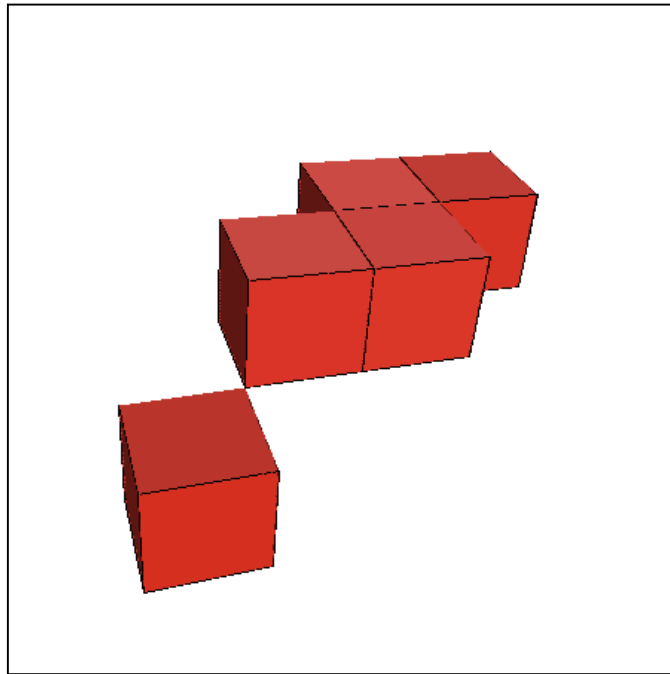
Recognition example



- $(0, 0, 0)$
- $(1, 1, 1)$
- $(2, 1, 1)$
- $(2, 2, 1)$
- $(3, 2, 1)$



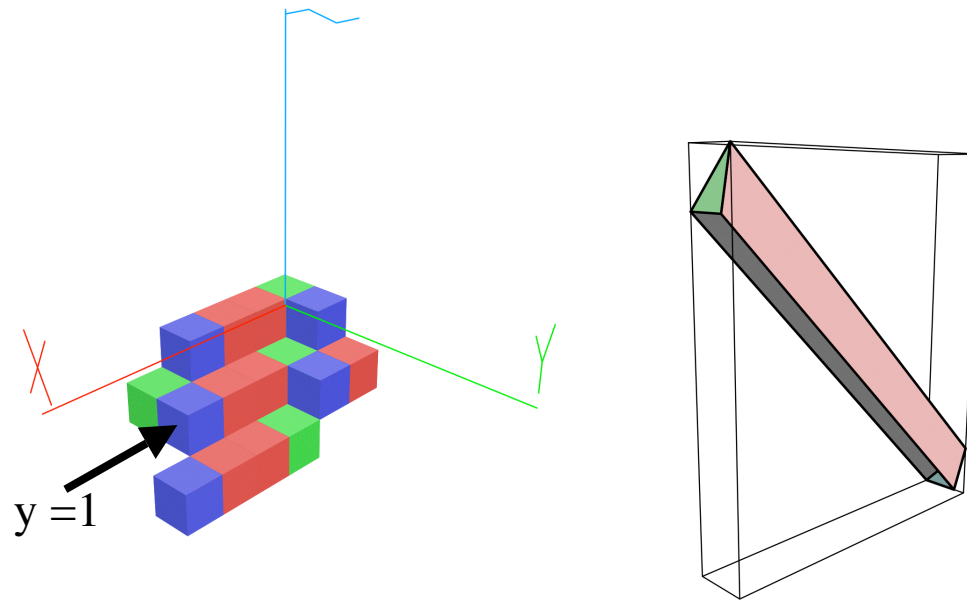
Recognition example



- $(0, 0, 0)$
- $(1, 1, 1)$
- $(2, 1, 1)$
- $(2, 2, 1)$
- $(3, 2, 1)$



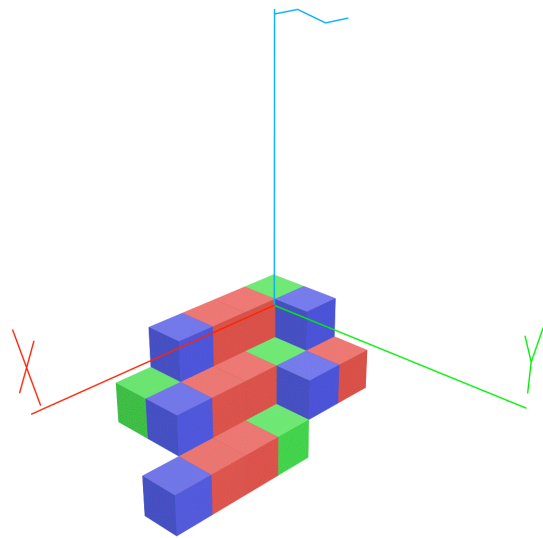
Preimage of a 3D straight line



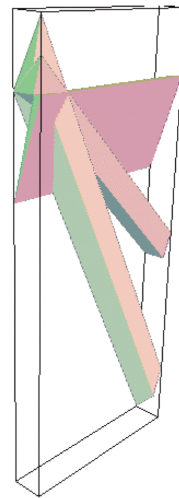
Preimage of the 3D digital straight line defined by $y = 1$ in the plane $P(1,3,4,0)$



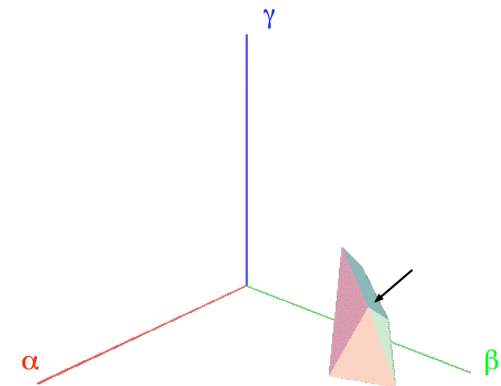
Preimage of a part of 3D naive plane



A piece of plane
 $P(1,3,4,0)$



Preimages of 3D
digital straight lines
in parameter space (a, b, g)



The preimage of the piece
of plane is the intersection
of the digital lines preimages.



Tricubes (1/2)

40 tricubes = 6 neutral + 17 pairs of symmetrical



[0]

(0,0,1)



[8]

(0,1,1)



[39]

(1,1,1)



[19]

(1,1,3)



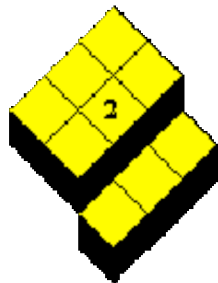
[13]

(1,2,3)



[32]

(2,2,3)



[2,6]

(0,1,2)



[24,30]

(1,1,2)



[15,35]

(1,2,2)

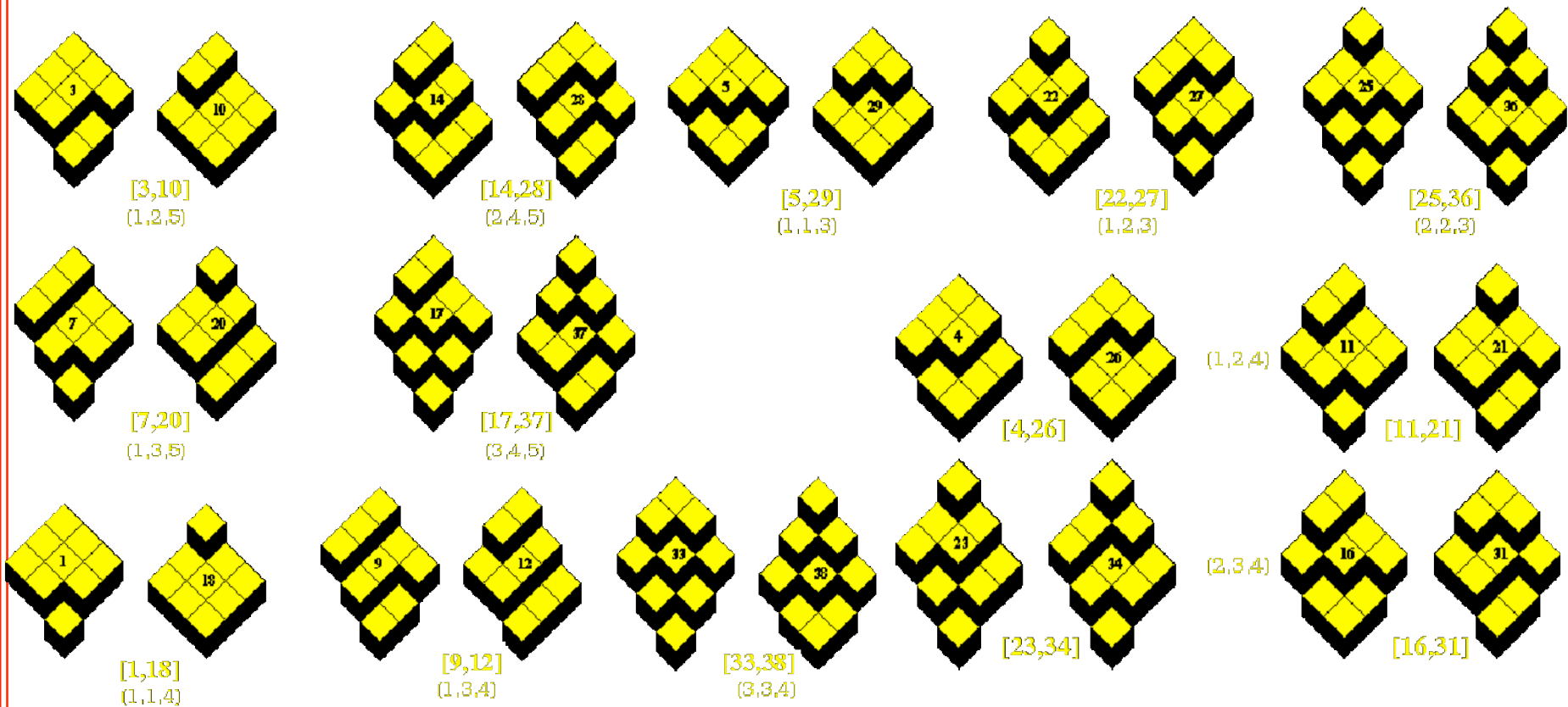


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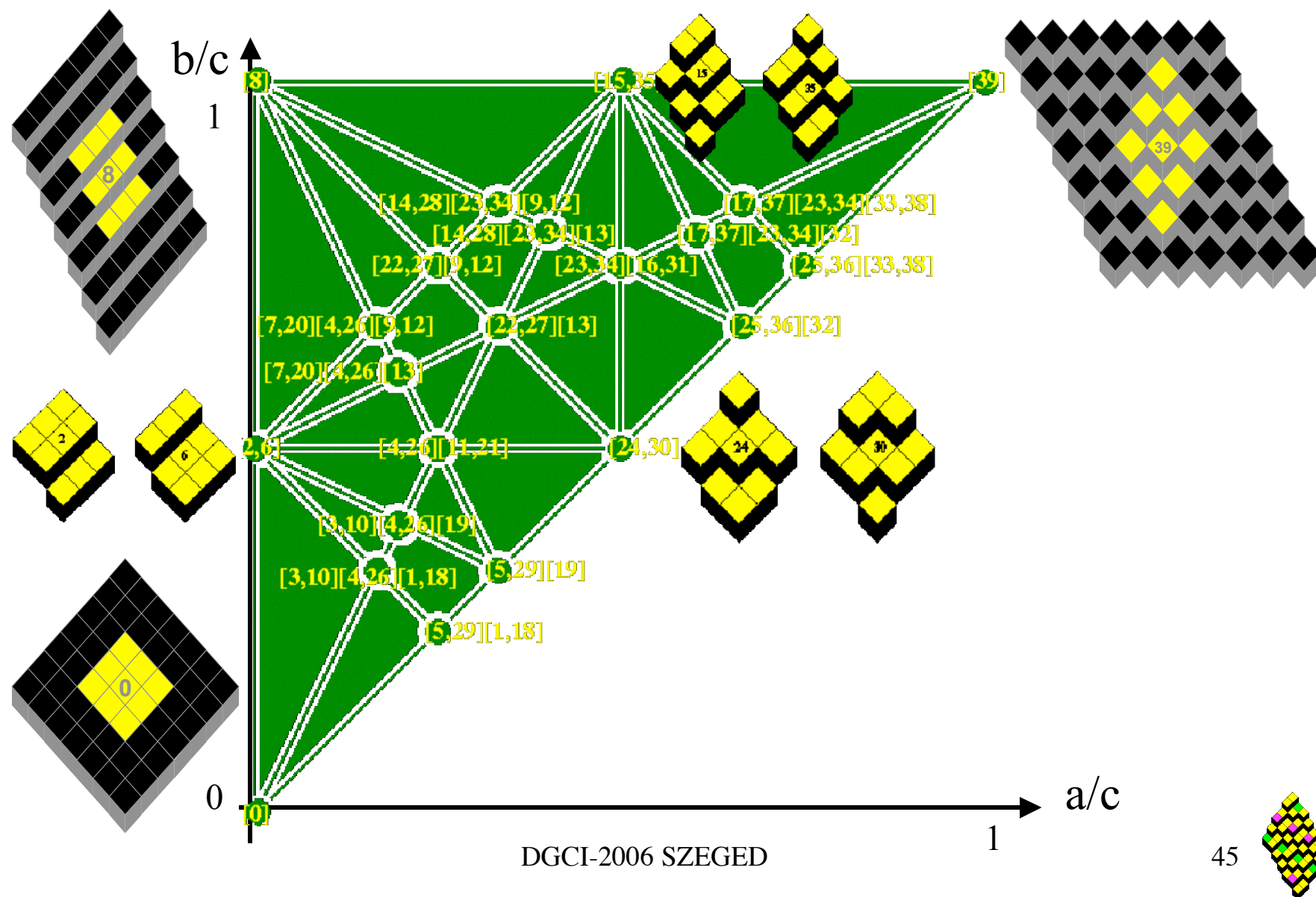
43



Tricubes (2/2)



Farey fan associated to tricubes



Tricubes - Properties

1. A naive plane contains no more than **9** different types of tricubes.
2. A **5x5** piece of naive plane centered on a leaning point contains **all** the tricubes used in the plane.
3. If the number of tricubes is :
 - even : all the tricubes are present by pairs
 - odd : one tricube is neutral and others are present by pairs



Tricubes - Properties

Lower Leaning Point

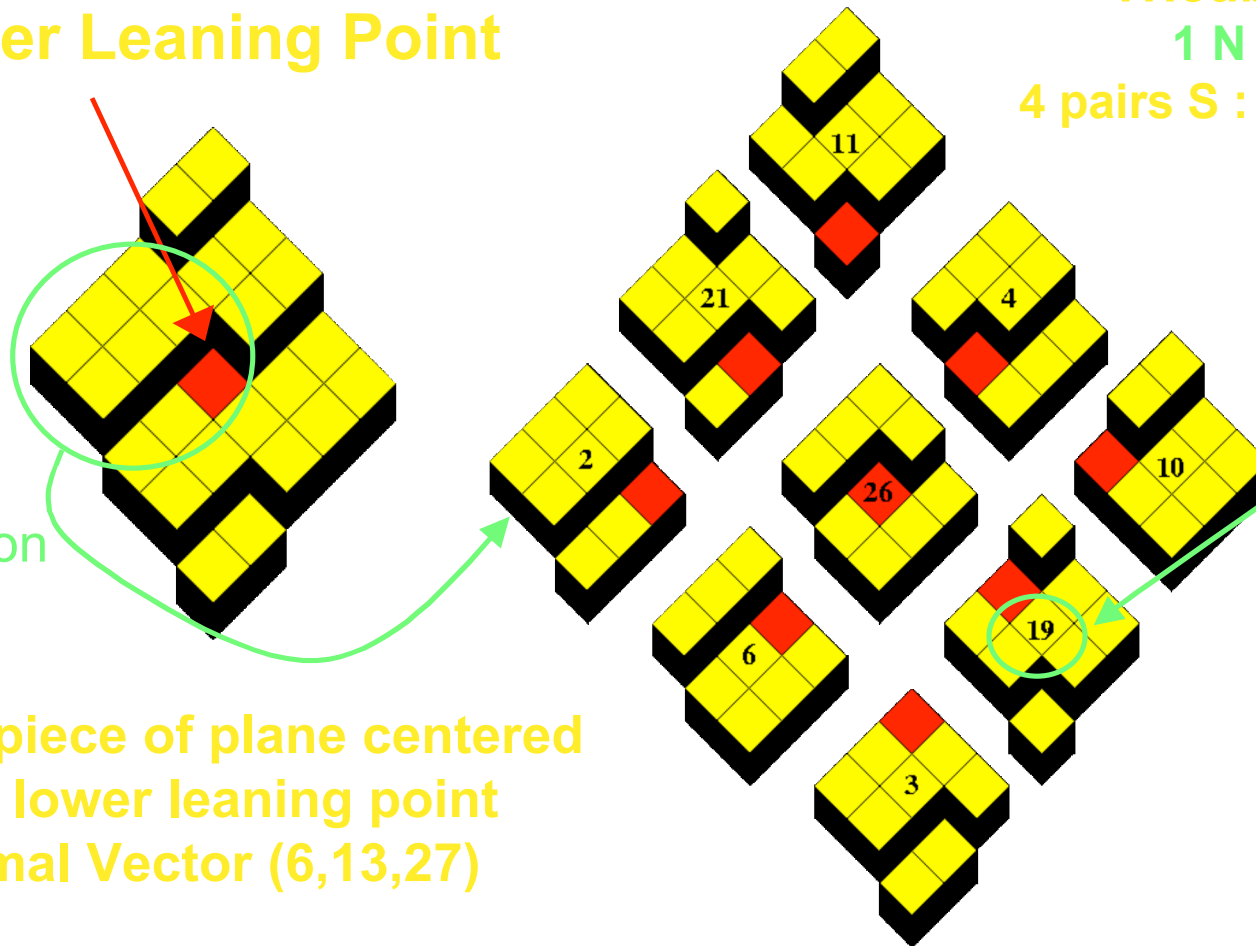
Exactly 9 types of
Tricubes :

1 N : 19

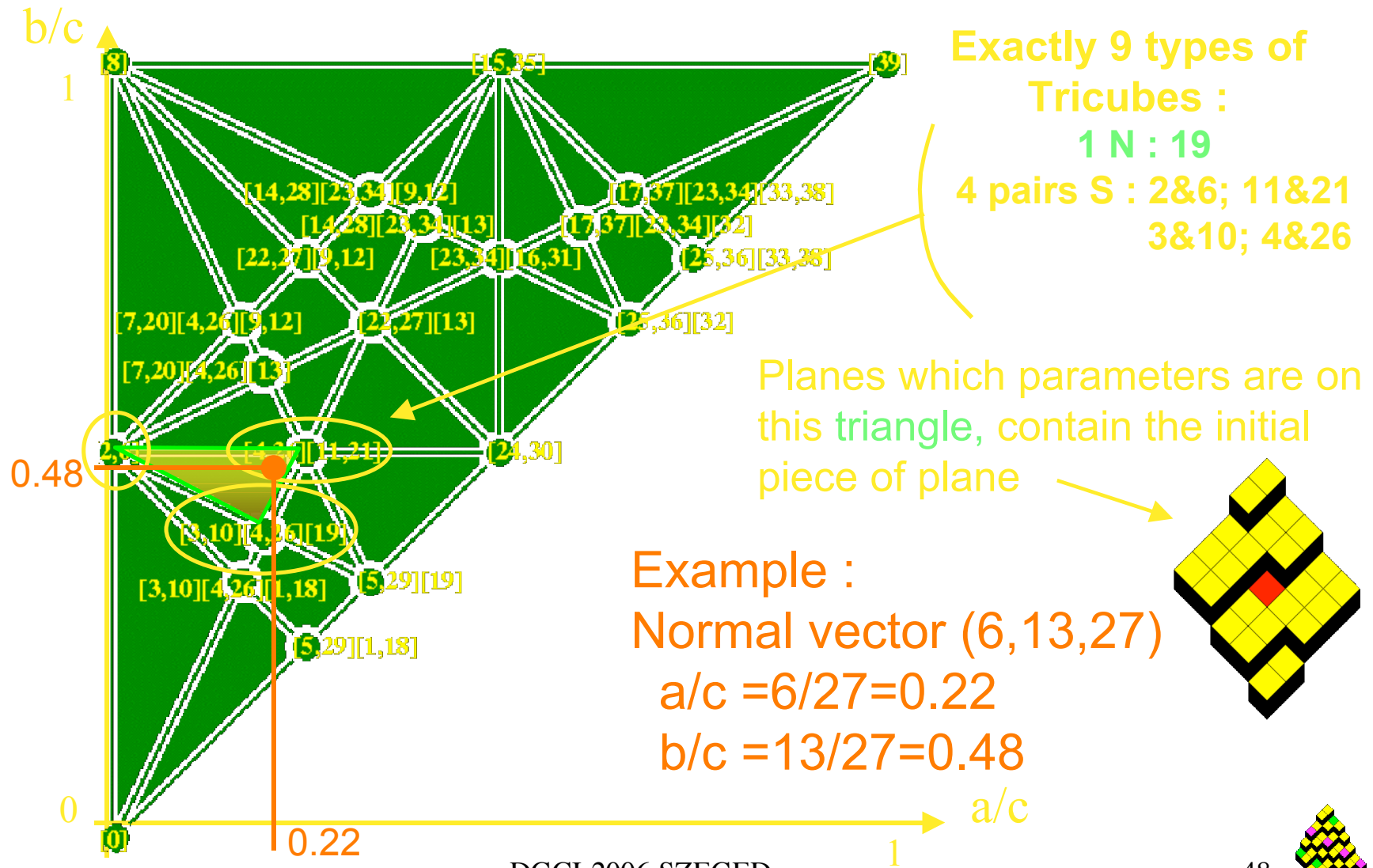
4 pairs S : 2&6; 11&21
3&10; 4&26

Tricubes
Decomposition

5x5 piece of plane centered
on a lower leaning point
Normal Vector (6,13,27)



Farey diagram associated to tricubes



Outline

1. Introduction
2. Hough transform and discrete geometry
3. Preimage of a pixel, a set of pixels, a connected straight line
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5. Extension in 3D
6. **Duality versus polarity**
7. Polygon enclosure and inclusion problems
8. Application to the convex approximation by interior.
9. Conclusion



Homogeneous representation

The homogeneous representation of a point $P = (x_1, x_2, \dots, x_d)$ in the d -dimensional Euclidean space is the point $(x_1, x_2, \dots, x_d, 1)$ in the projective space.

For any non-zero scalar λ , the homogeneous points $(\lambda x_1, \lambda x_2, \dots, \lambda x_d, \lambda)$ represent the same point in the Euclidean space.

This representation framework is convenient to obtain a matrix representation of both affine transformations and duality mapping.



Polarity

To define the polarity we consider the transformation:

$$\alpha = \beta B$$

with α and β vectors in $(d+1)$ homogeneous coordinate
and B a $(d+1) \times (d+1)$ matrix.

In 2D this transformation maps points to lines and lines to points.

This class of transformations preserves the incidence:

$$\alpha \in \beta, \text{ then } \beta B \in \alpha B$$

The matrix B is such that $\alpha B \alpha^T = 0$ which corresponds to equation of a conic in homogeneous coordinates.



Duality and polarity in 2D

Let us focus on polarity defined by the matrix:

$$\mathbf{B} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & -1 \end{bmatrix}$$

A point $(x_1, x_2, 1)$ is mapped to the dual line $x_1x + x_2y - z = 0$

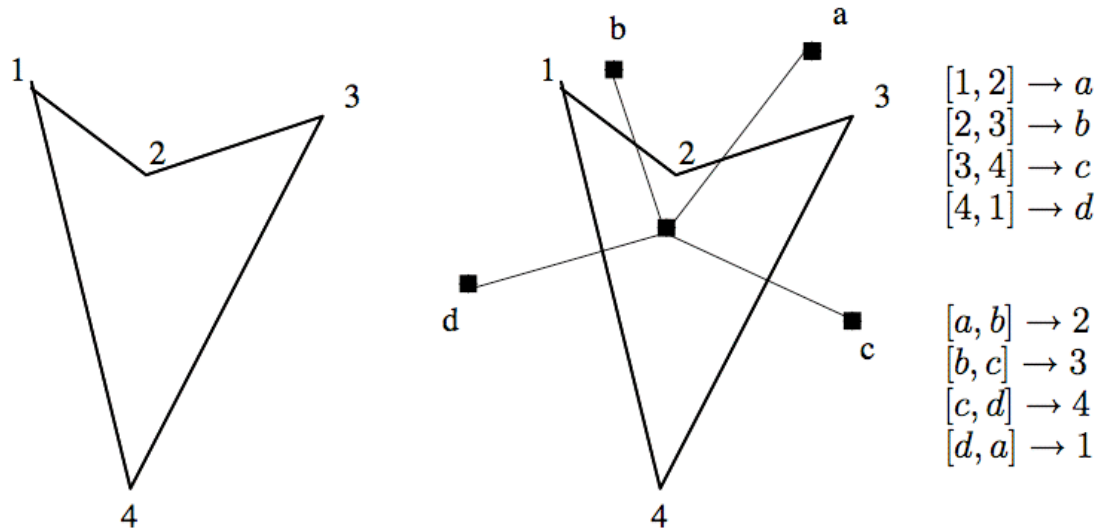
The conic defined by this mapping is a unit circle in 2D.

Given a point (or a line) α , we have:

$$\text{distance}(\alpha, O) \cdot \text{distance}(\text{dual}(\alpha), O) = 1$$



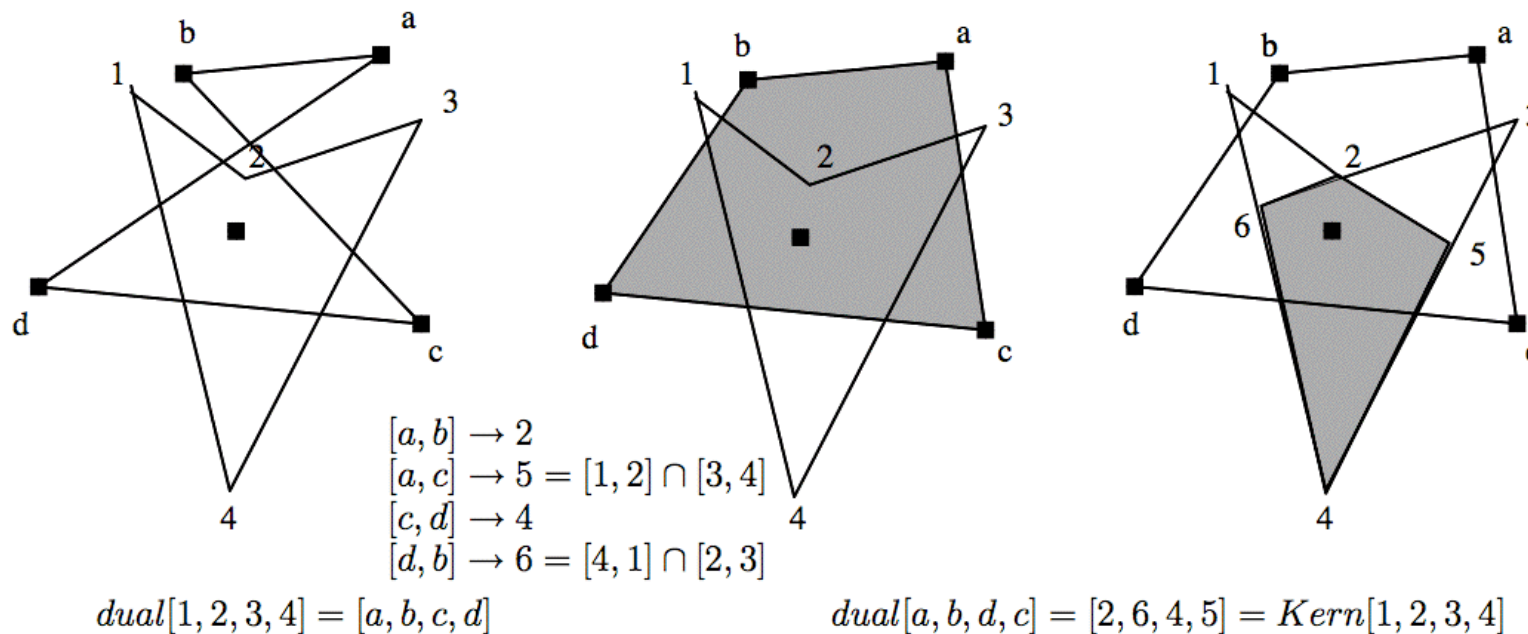
Polarity



Union and Intersection are “ dual ”



Polarity - Convexity - Kernel



The convex hull of the polar figure is the polar of the kernel.



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Polygon Enclosure problem

$$Enc(P, Q, \mu)$$

Given a general polygon P in class P ,

Find the μ -smallest polygon Q in class Q
with Q containing P .



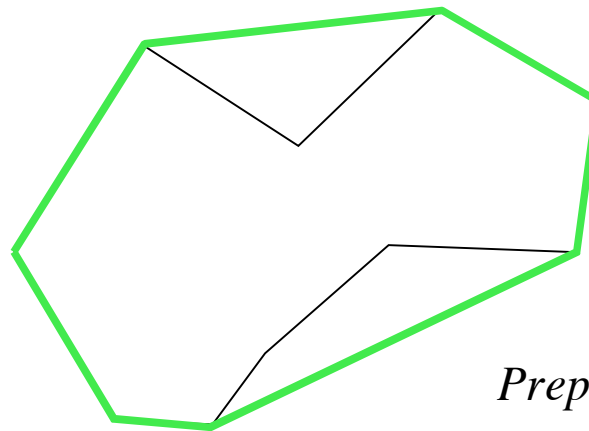
Polygon Enclosure problem

P is the class of n -sided simple polygons

Q is the class of **convex polygons**

μ is the area measure

$Enc(P, Q, \mu)$ is the problem of convex hull with complexity in $O(n \log n)$



Preparata and Shamos 1985



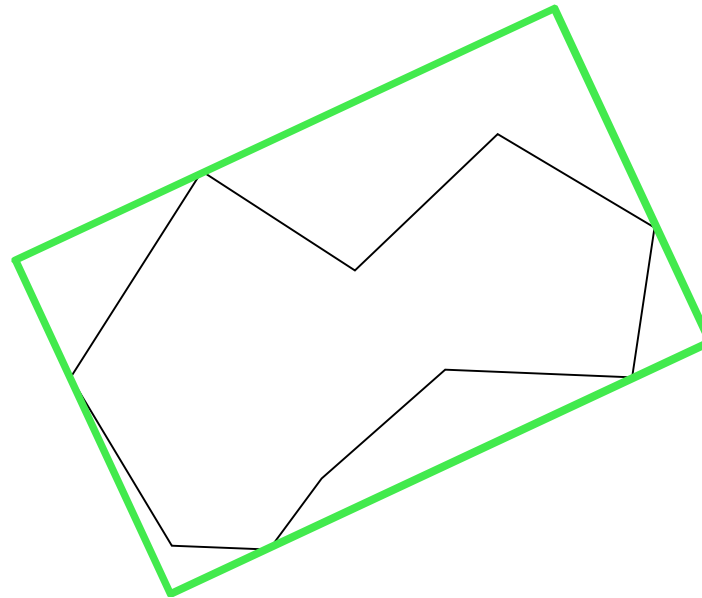
Polygon Enclosure problem

P is the class of n -sided simple polygons

Q is the class of rectangles

μ is the area measure

$Enc(P, Q, \mu)$ has a solution with complexity in $O(n)$ in 2D



Roth and al, 1997



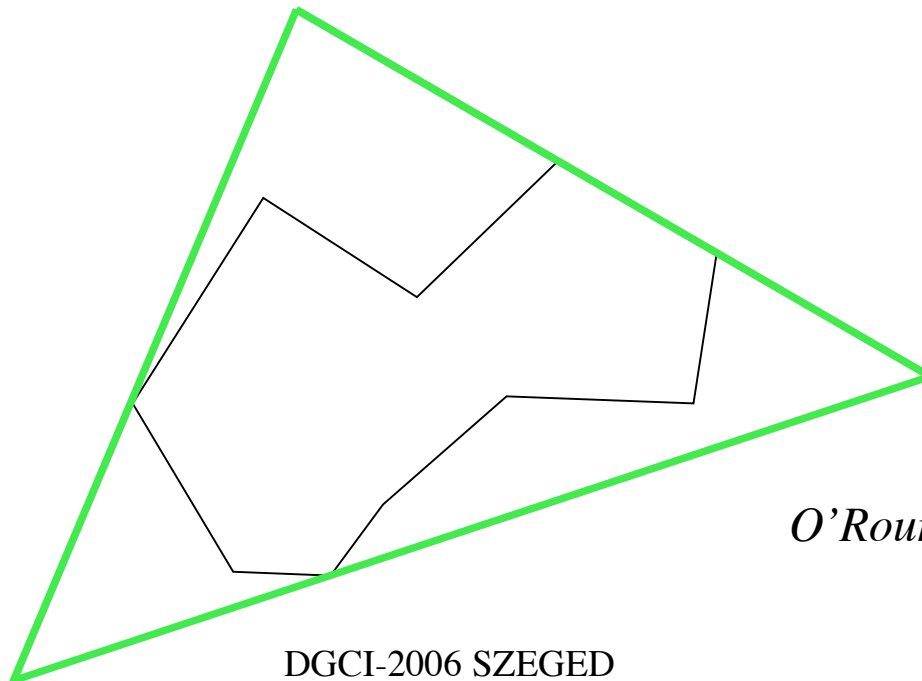
Polygon Enclosure problem

P is the class of n -sided simple polygons

Q is the class of triangles

μ is the area measure

$Enc(P, Q, \mu)$ has a solution with complexity in $O(n \log^2 n)$



O'Rourke and al 1986

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Polygon Inclusion problem

$$Inc(P, Q, \mu)$$

Given a general polygon P in class P ,

Find the μ -largest polygon Q in class Q

with Q contained in P .



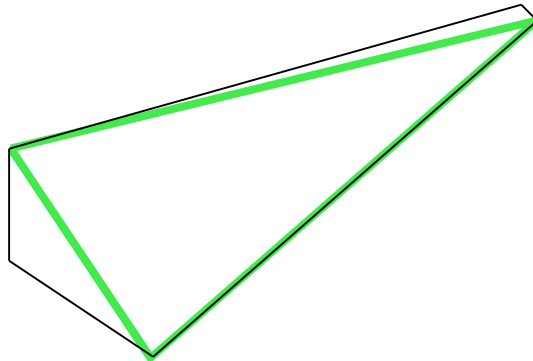
Polygon Inclusion problem

P is the class of n -sided convex polygons

Q is the class of **triangles**

μ is the area measure

$Inc(P, Q, \mu)$ has a solution with complexity in $O(n)$



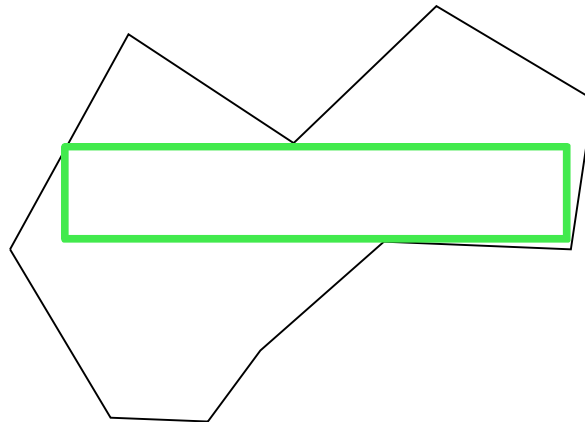
Polygon Inclusion problem

P is the class of n -sided simple polygons

Q is the class of **axis parallel rectangles**

μ is the area measure

$Inc(P, Q, \mu)$ has a solution with complexity in $O(n \log n)$



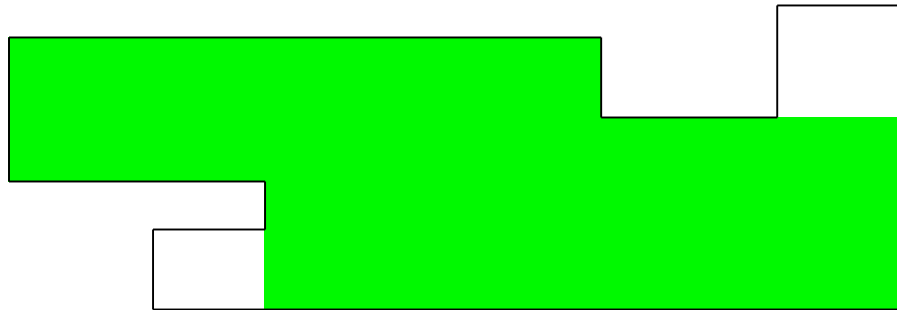
Daniels and al 1997



Polygon Inclusion problem

P is the class of n-sided orthogonal polygons
 Q is the class of **convex orthogonal polygons**
 μ is the area measure

$Inc(P, Q, \mu)$ has a solution with complexity in $O(n^2)$



Skull problem

Wood and Yap 1988



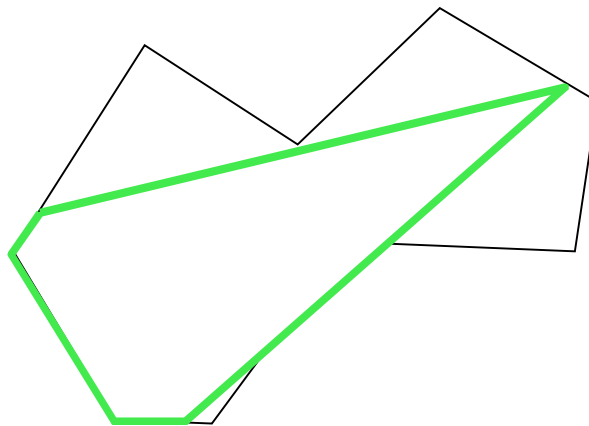
Polygon Inclusion problem

P is the class of n-sided simple polygons

Q is the class of **convex polygons**

μ is the area measure

$Inc(P, Q, \mu)$ has a solution with complexity in $O(n^7)$



Potato Peeling

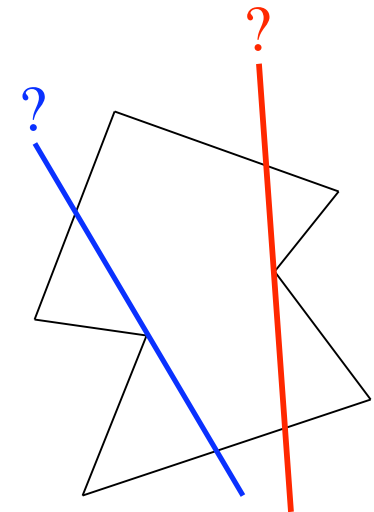
Chang and Yap 1986



Potato peeling (Chang & Yap 1986)

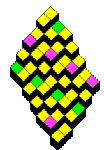
$Inc(P, Q, \mu)$

- Property: The Maximal Area Convex Subset (MACS) included in a polygon P is a polygon (Goodman 1981).
- So the MACS will be defined as the intersections between P and a set of half-planes.
- Half-planes will be associated to reflex vertices.



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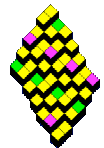


Heuristic approximation algorithm

$Inc(P, Q, \mu)$

To reduce the potato peeling algorithmic complexity, we propose an **approximation approach** based on kernel dilation.

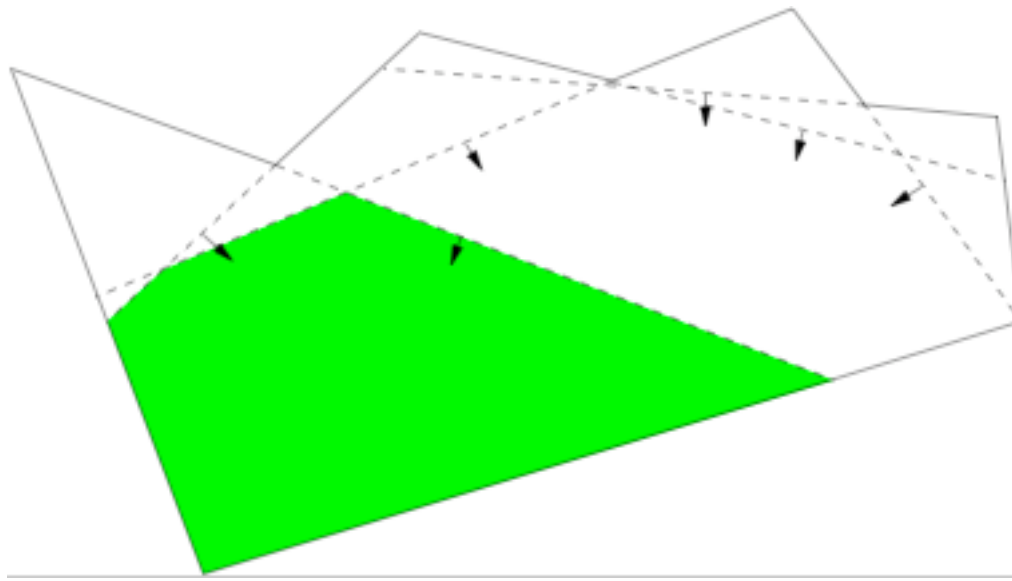
Chassery & Coeurjolly, ISMM, 2005



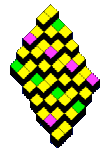
Star-shaped polygon and kernel

Definition: P is a star-shaped polygon if there exists a point q in P such that the segment qv_i lies inside P for all vertices v_i of P . The set of such q points is the Kernel of P .

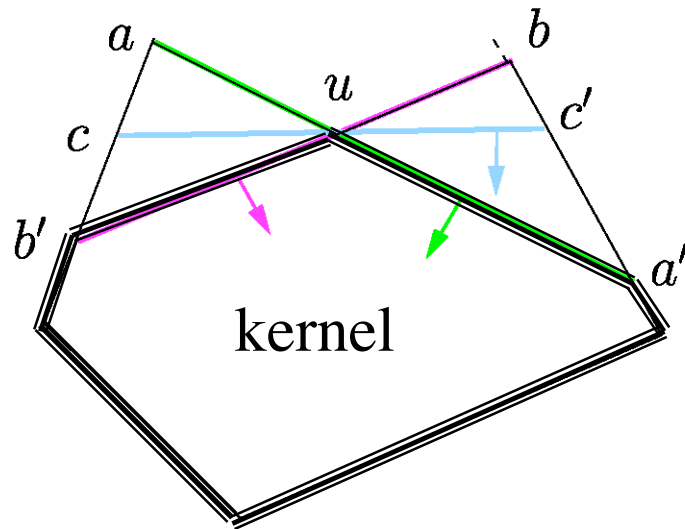
Proposition: the Kernel of P is given by intersections between P and the half-planes defined by all « extremal chords » associated to all reflex vertices.



Kernel = dual of the convex hull of the polar polygon P



Limitation to star-shaped polygon



The kernel of P is a convex subset of the MACS.

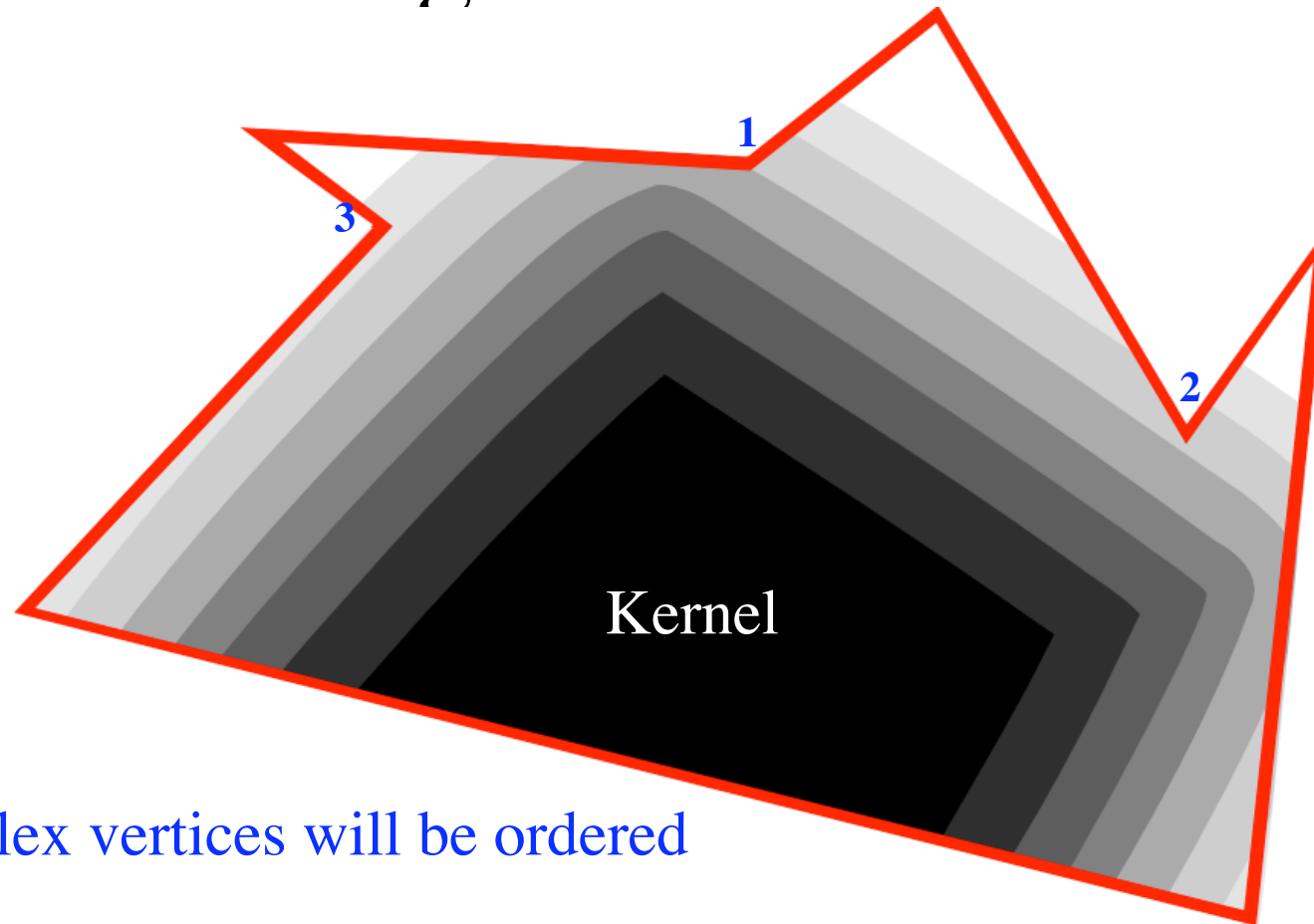
How to deform the kernel to approximate the MACS ?

The deformation will be an Euclidean dilation of the kernel.

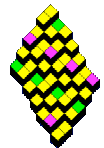


Level curves issued from kernel

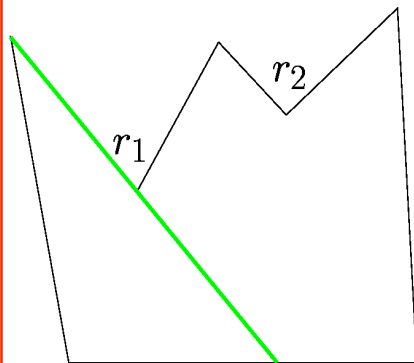
Ordering the reflex vertices



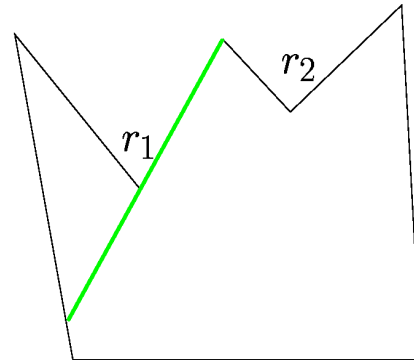
The reflex vertices will be ordered
by the dilation wavefront.



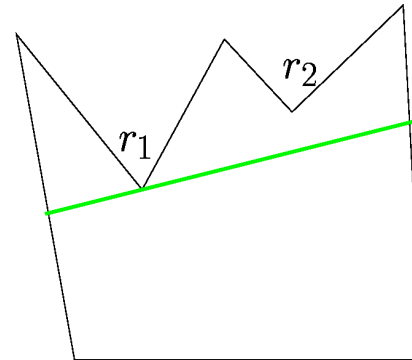
Definition: A Chord is a Maximal segment fully contained in P.



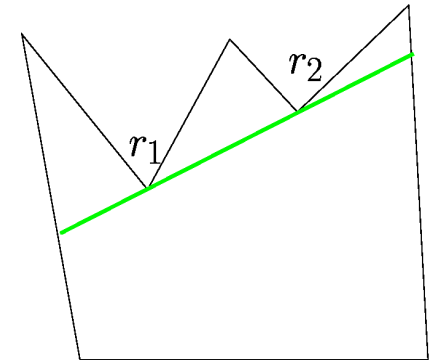
(a)



(a)



(b)



(c)

When a reflex vertex is analyzed, we have different possibilities for chords:

The chord may be an **extremal chord**.

(a)

The chord may be a **single-pivot chord** (its slope is tangent to the wavefront).

(b)

The chord may be a **double-pivot chord** between two successive reflex vertices.

(c)

From these possibilities, we choose the chord that maximizes the area of the resulting polygon.

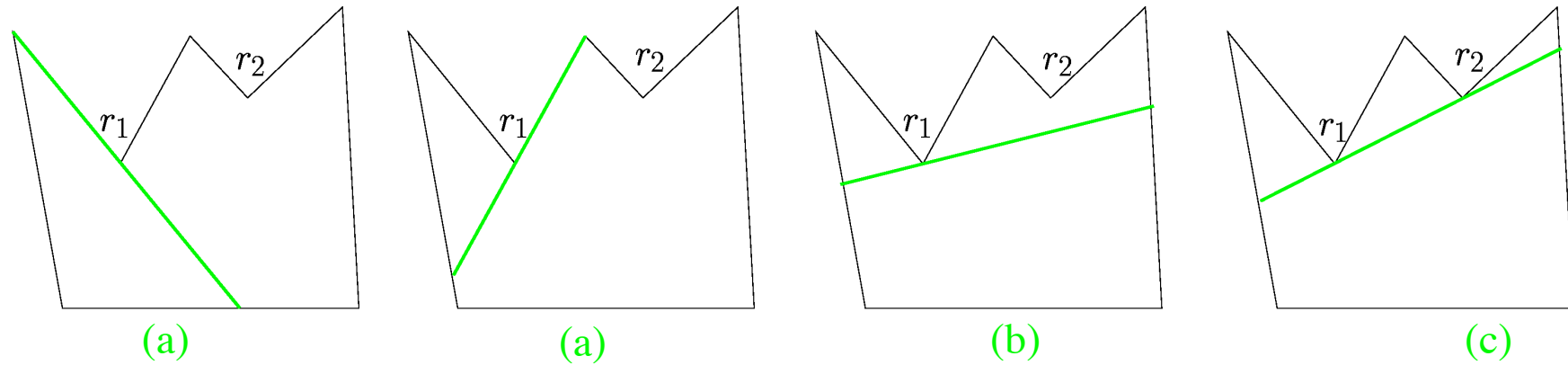


Approximated MACS algorithm

1. Compute the kernel of P
2. Compute the ordered list O of reflex vertices
3. Extract the first point r_1 in O
4. While O is not empty do
 5. Extract the first point r_2 in O
 6. Choose the best chord that maximizes the resulting polygon area with the chords (r_1, r_2)
 7. Modify the polygon P accordingly
 8. Update the list O removing reflex points excluded by the chord
 9. $r_1 \leftarrow r_2$
10. End while

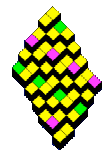
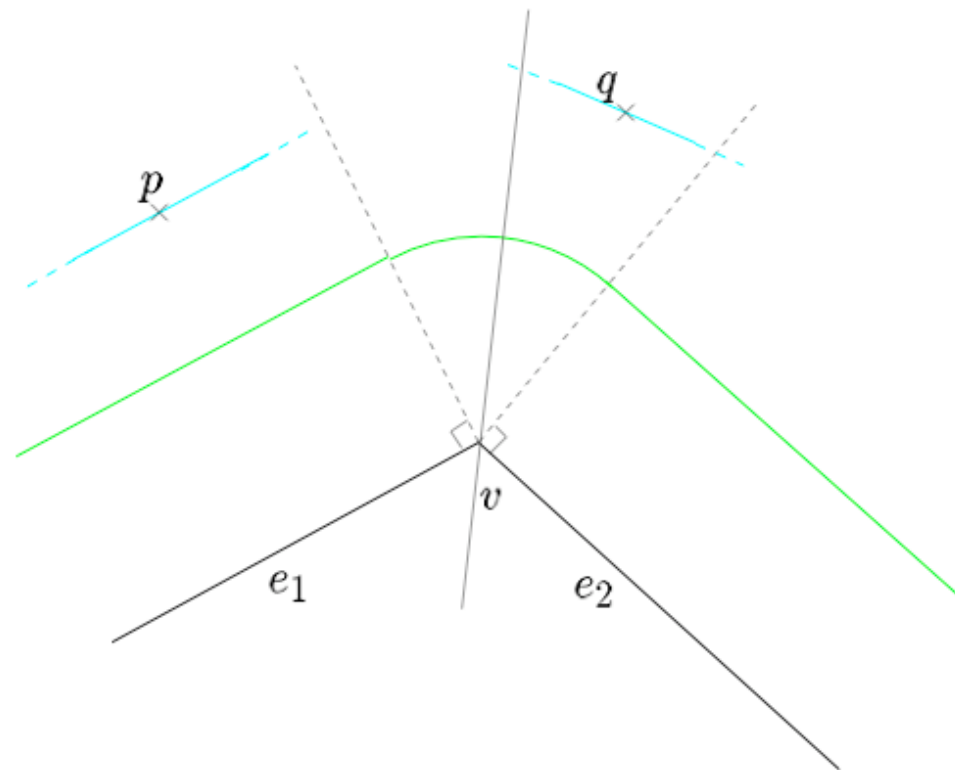


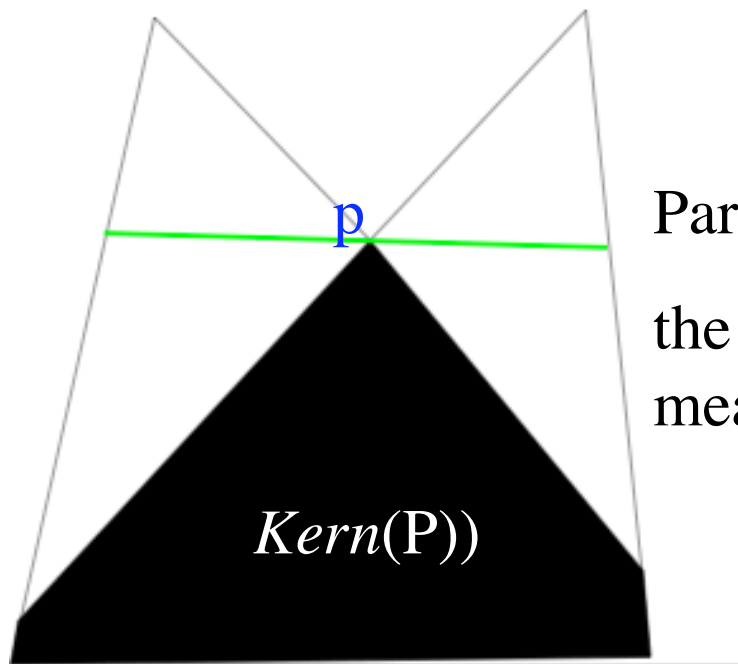
Slope computation



The slope computation is direct for **extremal chord (a)** and **double-pivot chord (c)**.

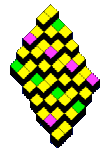
For a single-pivot chord :
the chord is tangent to the wavefront issued from the Kernel.



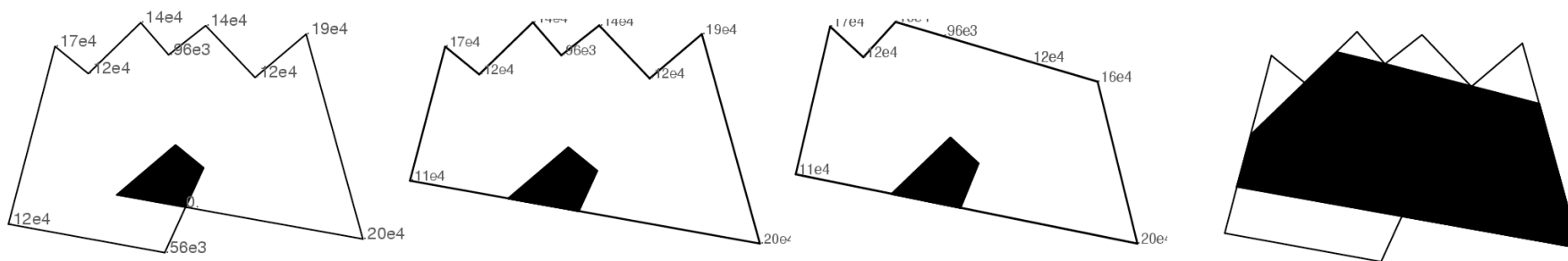


Particular case where $d(p, Kern(P)) = 0$

the single-pivot chord slope is obtained as the mean between the two maximal chord slopes.



Intermediate steps of the approximated MACS algorithm



Kernel

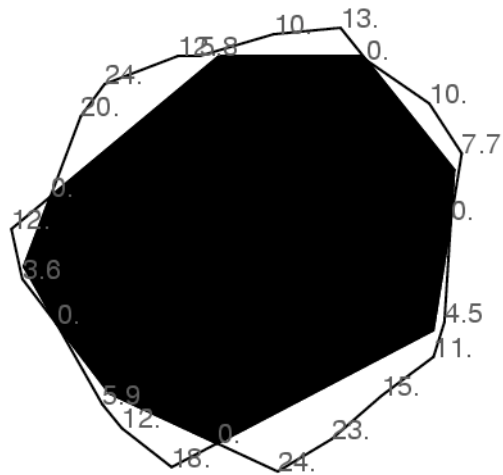
1st cut :
extremal chord

2nd cut :
double-pivot chord

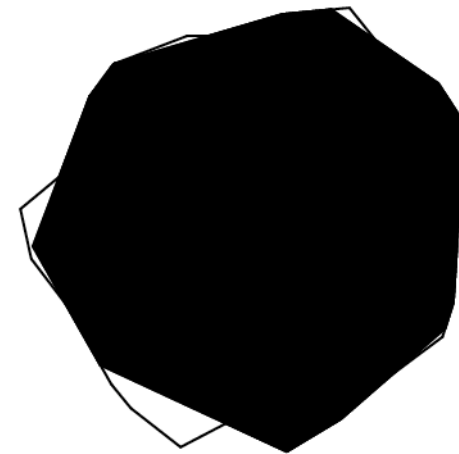
3rd cut :
extremal chord



Results

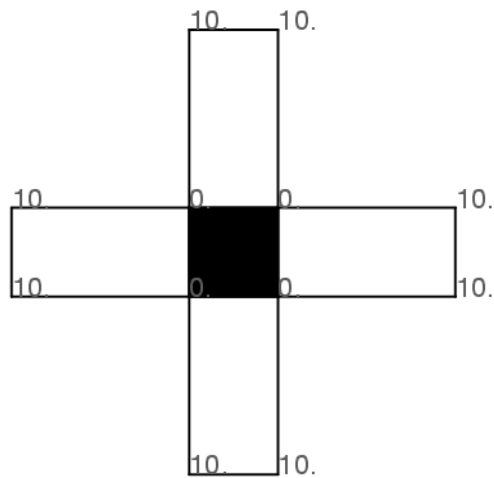


Initial polygon and its kernel



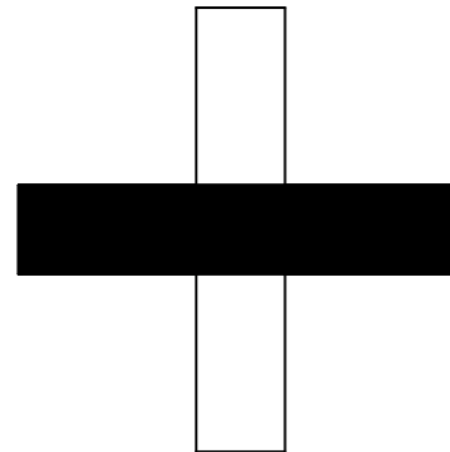
Approximated MACS





Initial polygon and its kernel

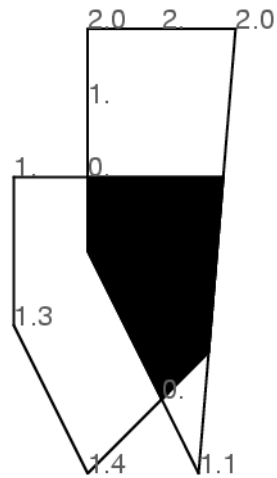
Results



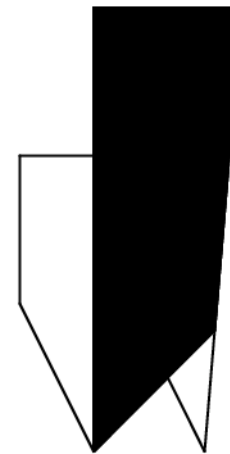
Approximated MACS



Results



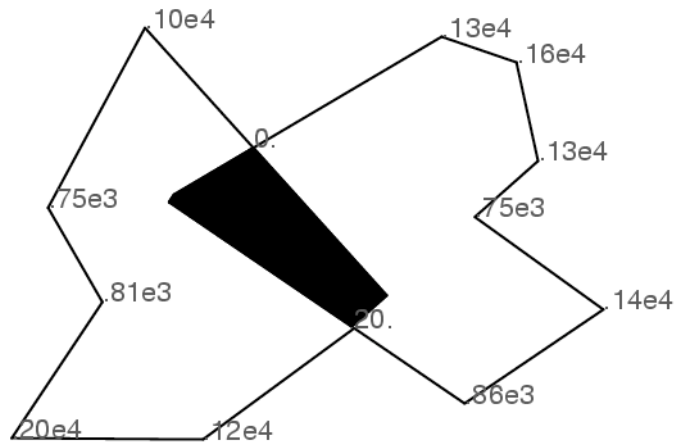
Initial polygon and its kernel



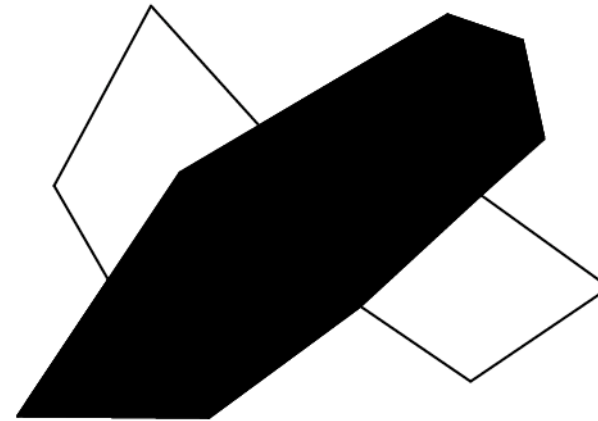
Approximated MACS



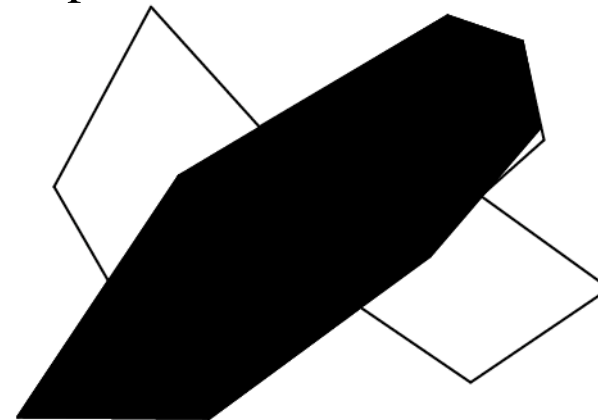
Comparison between Optimal MACS and Approximated MACS



Initial polygon and its kernel



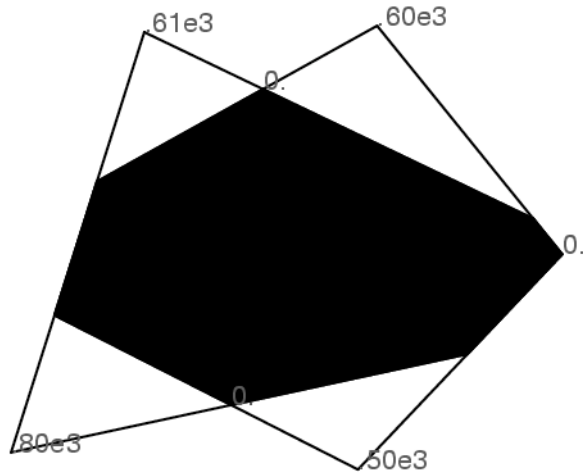
Optimal MACS



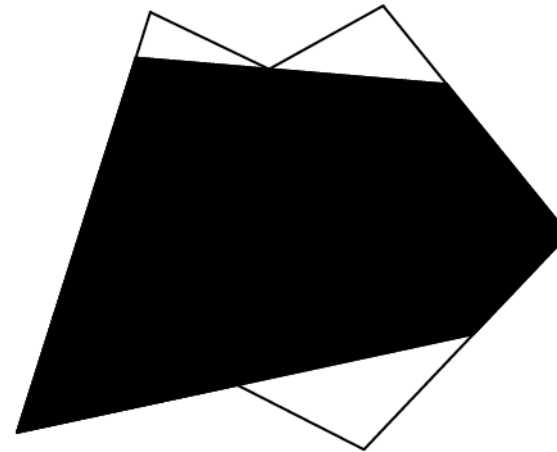
Approximated MACS



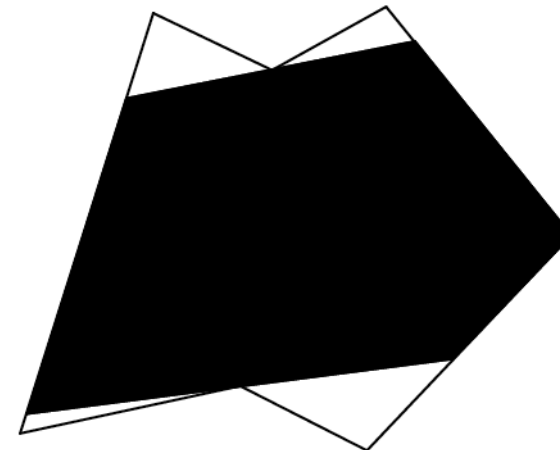
Comparison between Optimal MACS and Approximated MACS



Initial polygon and its kernel



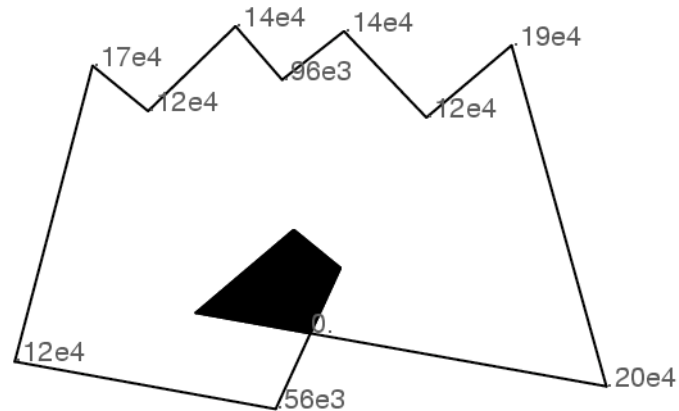
Optimal MACS



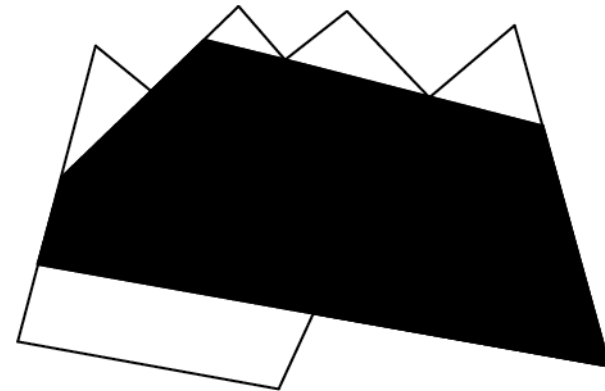
Approximated MACS



Comparison between Optimal MACS and Approximated MACS



Initial polygon and its kernel



Optimal or Approximated MACS



Computational cost

n : number of vertices of P

k : number of reflex vertices

Complexity in $O(n + k \log k)$

1. Compute the kernel of P
2. Compute the ordered list O of reflex vertices
3. Extract the first point r_1 in O
4. While O is not empty do
 5. Extract the first point r_2 in O
 6. Choose the best chord that maximizes the resulting polygon area with the chords (r_1, r_2)
 7. Modify the polygon P accordingly
 8. Update the list O removing reflex points excluded by the chord
 9. $r_1 \leftarrow r_2$
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- $O(n)$
- $O(k)$
- $O(k \log k)$
- $O(k)$
- $O(k)$



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Conclusion

Two approaches for duality:

- parametric transform
- geometric transform

Duality offers new representation space for

- characterization,
- recognition,
- optimization.

Duality versus polarity has been illustrated.

Such work may be extended.



Tricubes - Properties

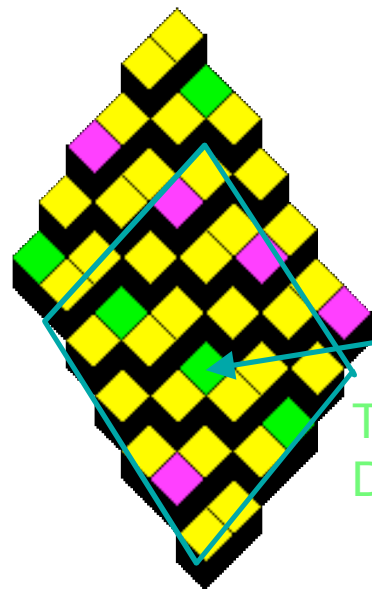
7 types of Tricubes :

1 N : 32

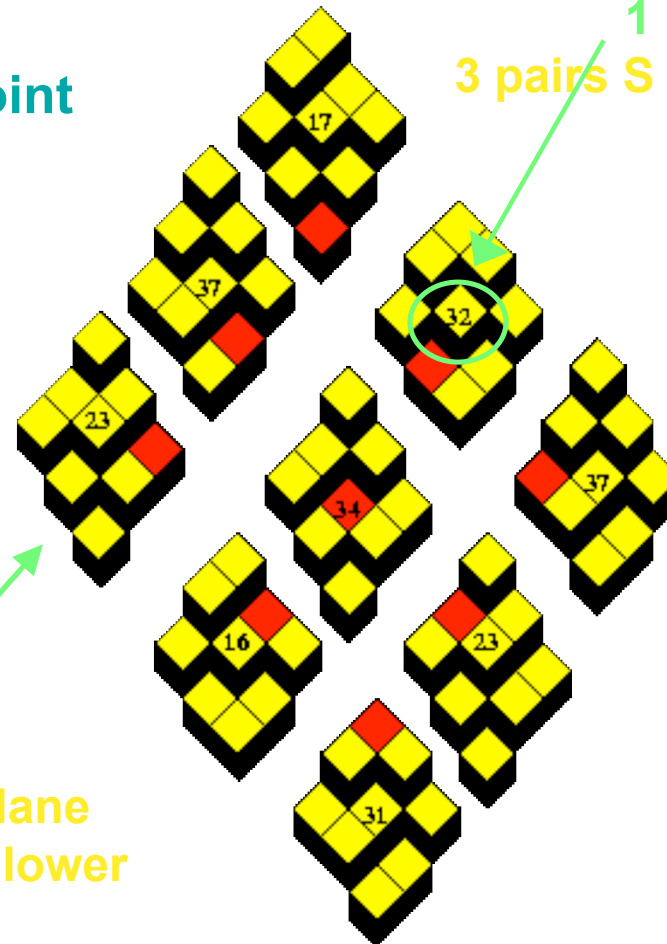
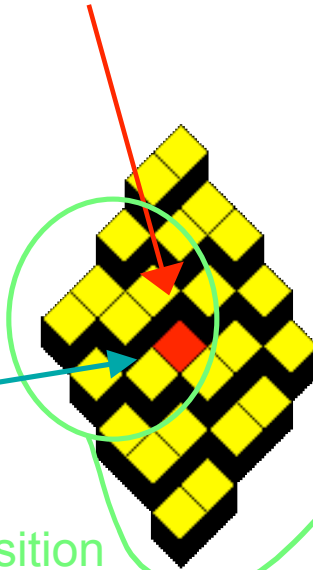
3 pairs S : 23&34
17&37
16&31

Piece of plane
Normal Vector (5,7,9)

Lower Leaning Point



Tricubes
Decomposition

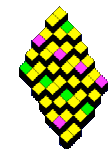
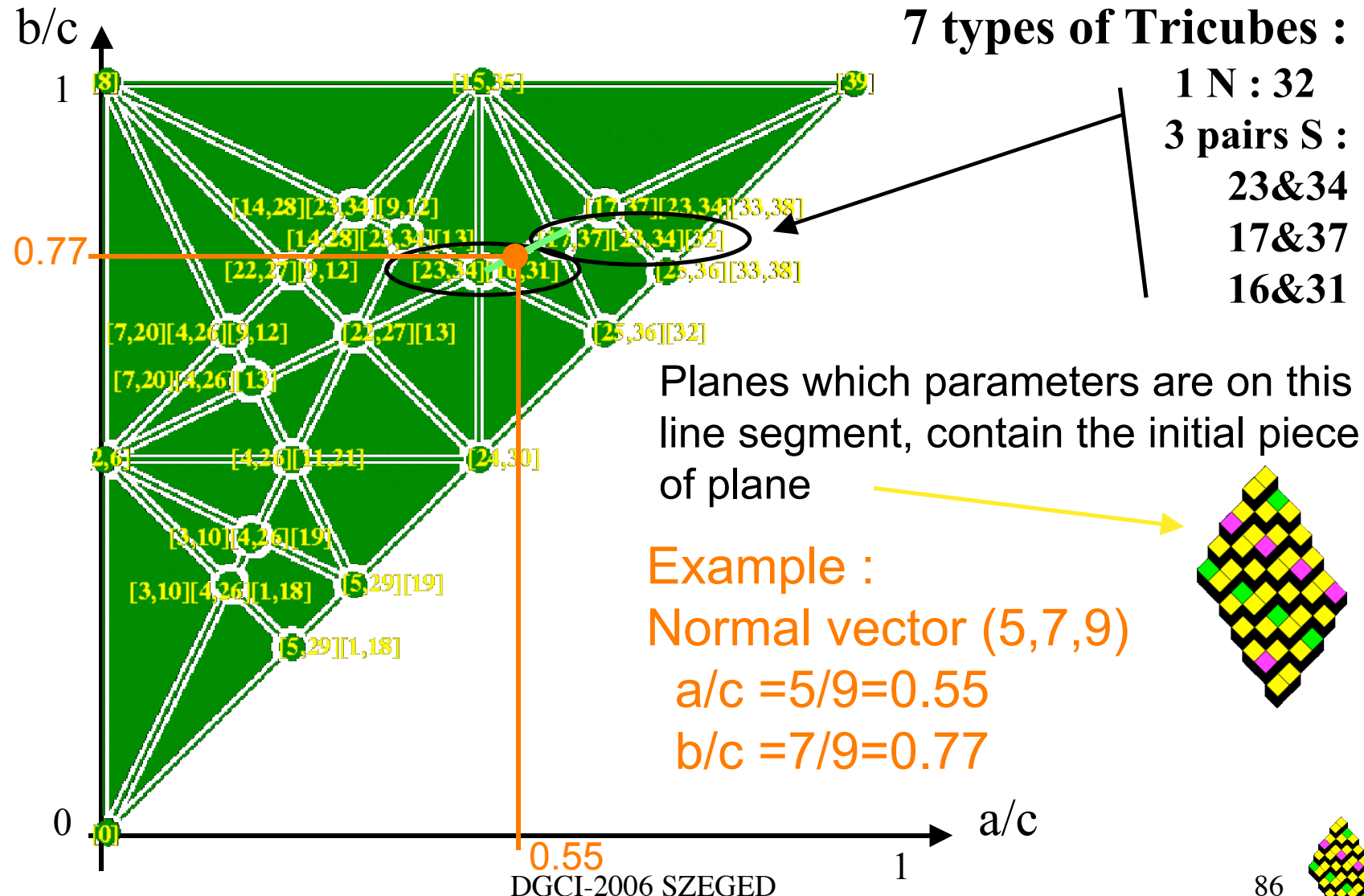


Lower leaning points
Upper leaning points

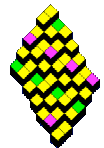
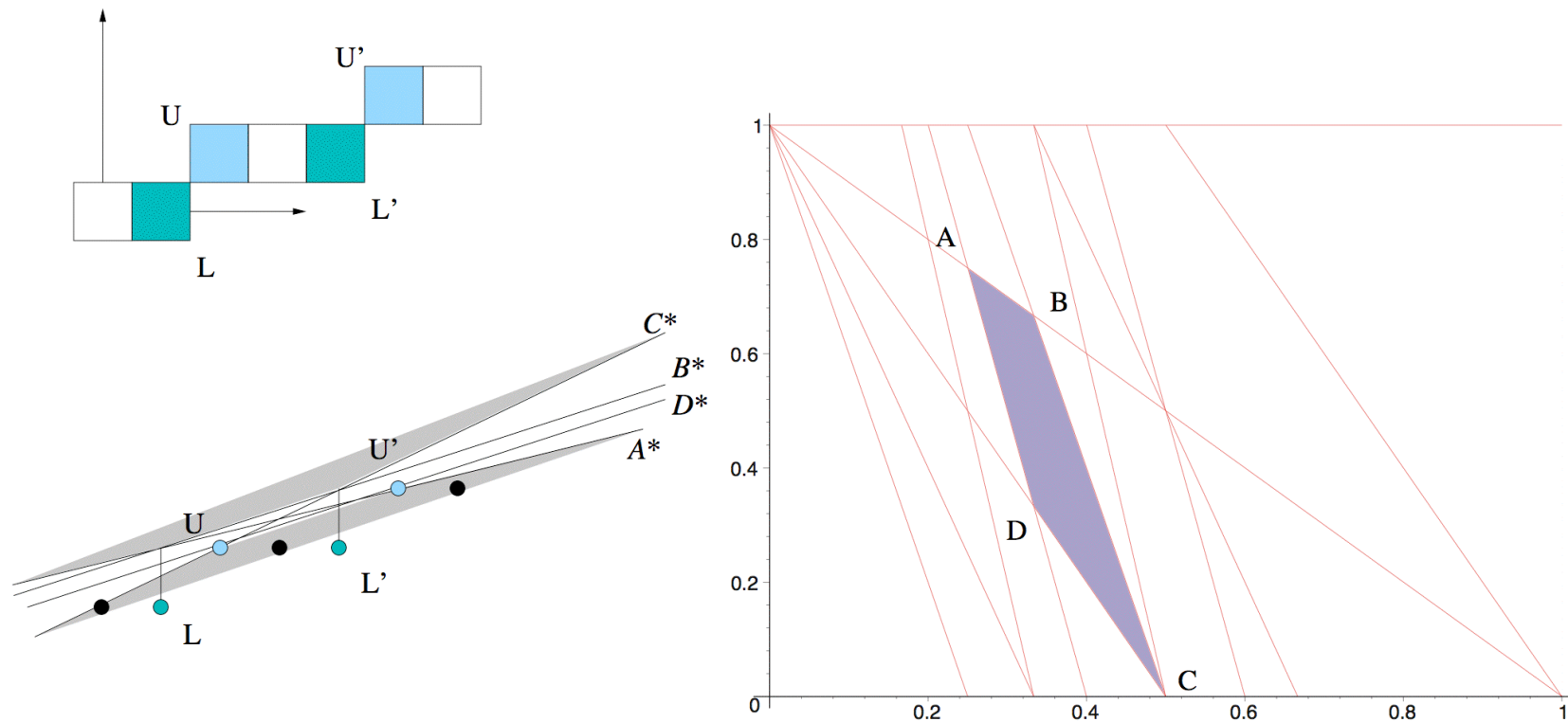
5x5 piece of plane
centered on a lower
leaning point



Farey diagram associated to tricubes

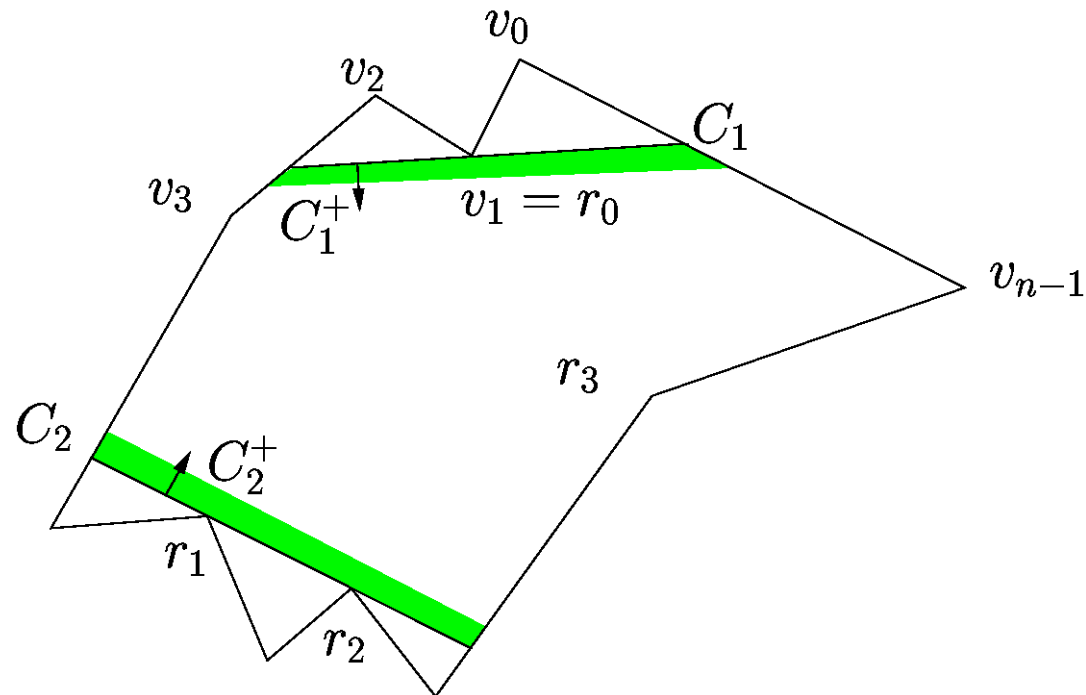


Relation between digital straight line segment and duality



Optimal solution is based on finding constraints on different reflex vertices:

- in case of one isolated reflex vertex --> an infinity of solutions (slope of the constraint)
- in case of consecutive reflex vertices --> one solution



Optimal solution is obtained by intersection of all the constraints.

